Pooling or Sampling: Collective Dynamics for Electrical Flow Estimation

Emanuele Natale
joint work with L. Becchetti[†] and V. Bonifaci[*]

[†] Sapienza Università di Roma

[*] National Research Council of Italy

AG1 Mittagsseminar

24 May 2018

Electrical Networks for Optimization

Computation of currents and voltages in resistive electrical network is a crucial primitive in many optimization algorithms

- Maximum flow
  - Christiano, Kelner, Madry, Spielman and Teng, STOC'11
  - Lee, Rao and Srivastava, STOC'13
- Network sparsification
  - Spielman and Srivastava, SIAM J. of Comp. 2011
- Generating spanning trees
  - Kelner and Madry, FOCS'09

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and as model of biological computation

- Physarum polycephalum
- Ants

implicitly solving electrical flow while forming food-transportation networks

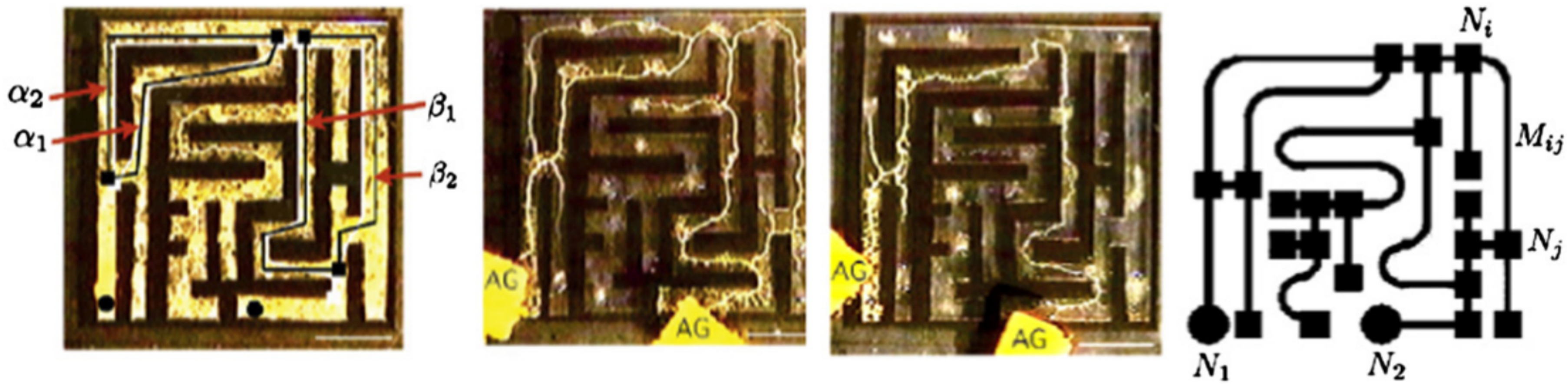
The Slime Mold *Physarum Polycephalum*

Many nice videos on Youtube:

https://bit.ly/1T5cSSY

The Slime Mold *Physarum Polycephalum*

Nakagaki, Yamada and Toth, Nature 2000
Tero, Kobayashi and Nakagaki J. of Theo. Bio. 2007



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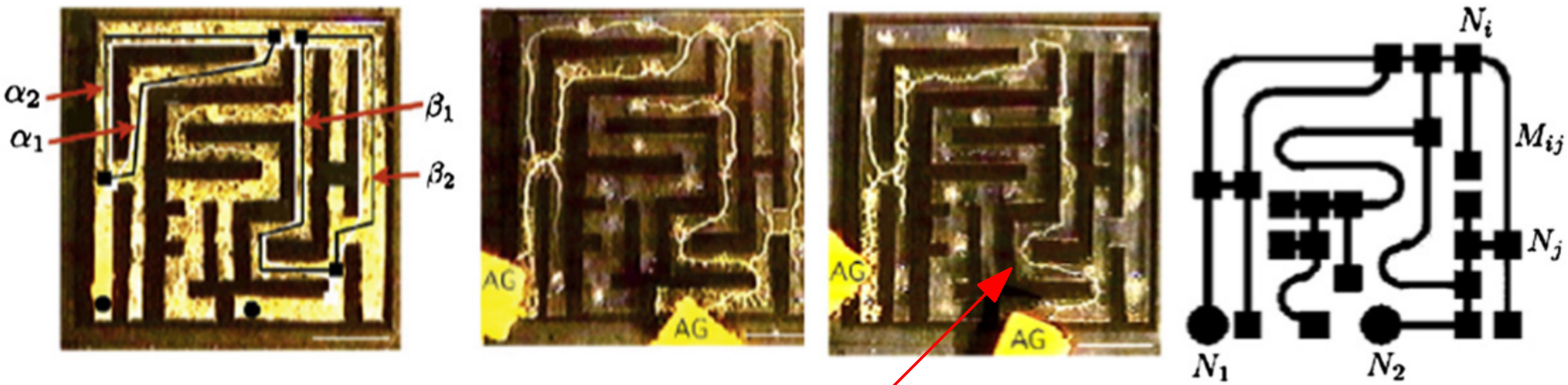

Physarum polycephalum builds *tubes* to transport food.
The amount of food flowing in the tube determines its growth or deterioration.

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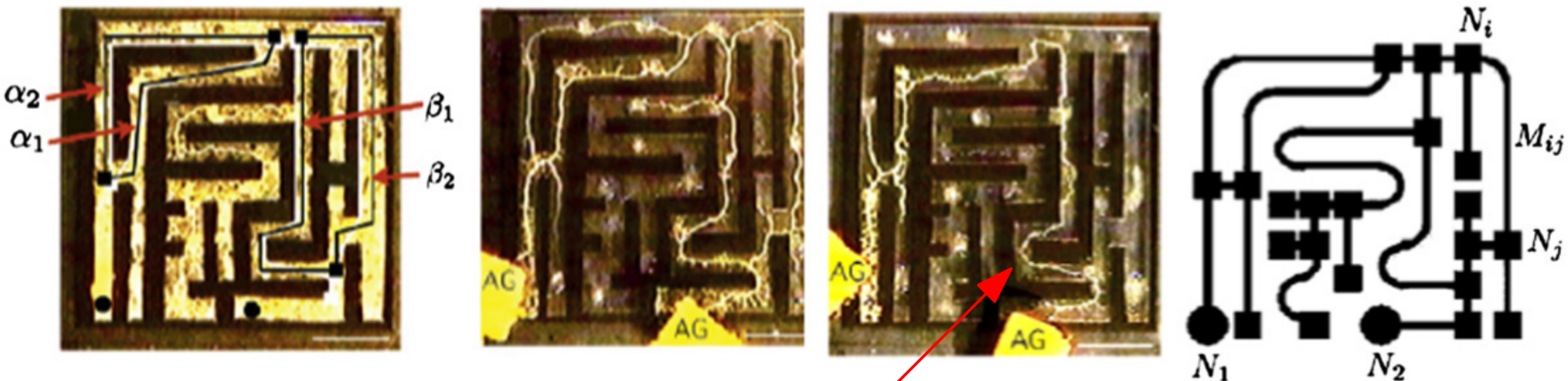

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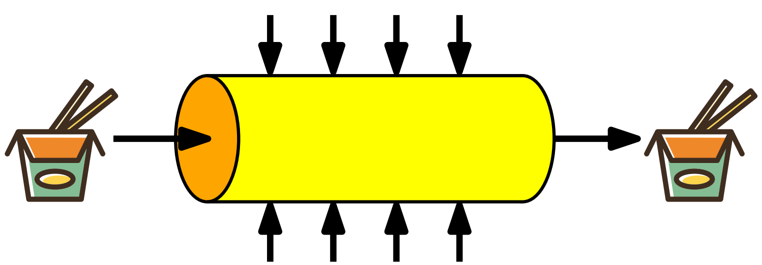

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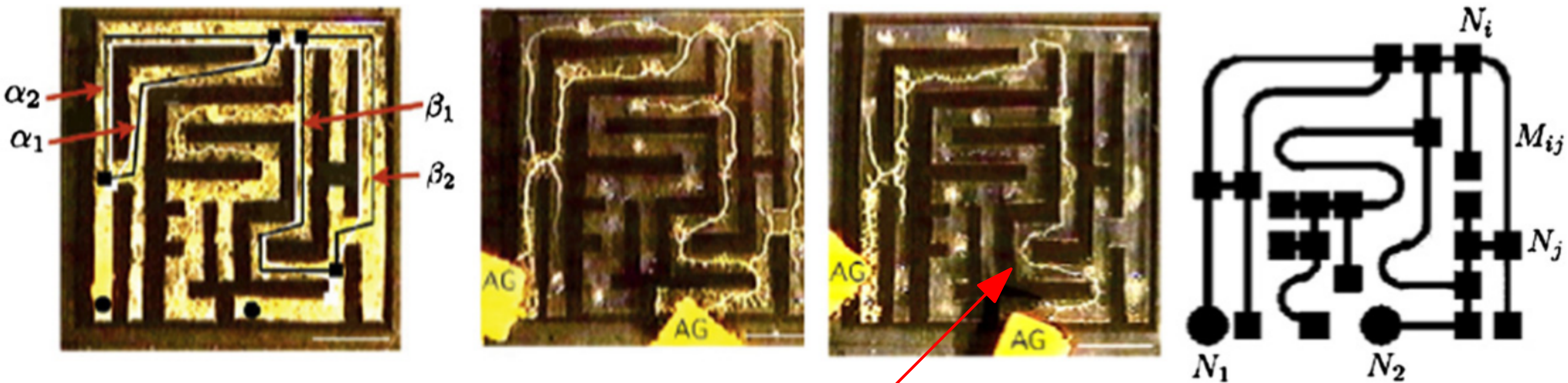

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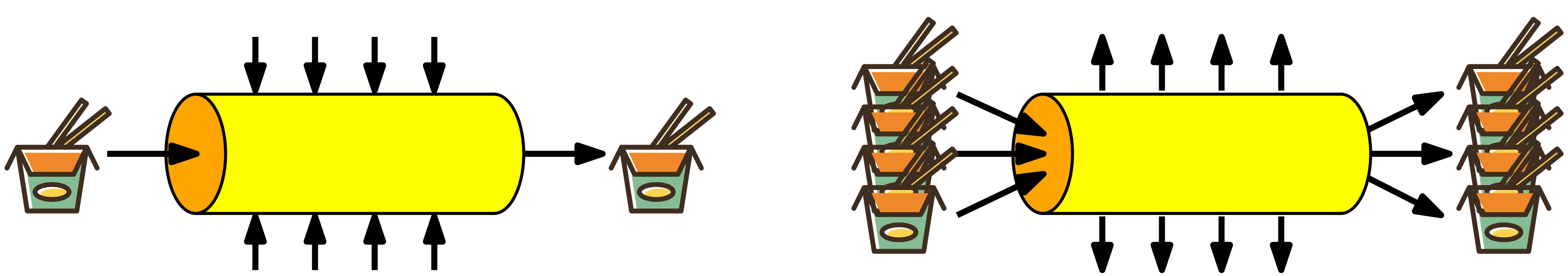

The Slime Mold *Physarum Polycephalum*

For each edge e and node u

- ℓ_e length
- x_e thickness
- q_e food flow
- $r_e = \ell_e / x_e$ resistance to flow

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- Flows relates to pressures by
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Dynamics: $\dot{x}_e = |q_e| - x_e$

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Is the model *good*?
See **S**lime **M**old **G**raph **R**epository, Dirnberger et al.

Electrical Networks as Biological Models?...

hydraulic	electric
volume V [m^3]	charge q [C]
pressure p [Pa=J/m^3=N/m^2]	potential ϕ [V=J/C=W/A]
Volumetric flow rate Φ_V [m^3/s]	current I [A=C/s]
velocity v [m/s]	current density j [C/(m^2·s) = A/m^2]
Poiseuille's law $\Phi_V = \frac{\pi r^4}{8\eta}\frac{\Delta p^\star}{\ell}$	Ohm's law $j = -\sigma\nabla\phi$

(stolen from Wikipedia/Hydraulic_analogy)

From Slime Molds to Ants

electric network	*Physarum*	ant trails
length in space	length in space	length in space
potential/voltage	amount of nutrient	number of ants
current	flow of nutrient	flow of ants
conductivity	thickness of tube	pheromone concentration
capacitance	transport efficiency	total pheromone density
reinforcement intensity	tube expansion rate	pheromone drop rate
conductivity decrease rate	tube decay rate	evaporation rate

For each edge e and node u

- ℓ_e length
- x_e conductivity
- q_e current
- $r_e = \ell_e / x_e$ resistance

- $b(u)$ 1 on source, -1 on sink, 0 o/w
- Kirchhoff current law: $\sum_{v \sim u} q_{(u,v)} = b(u)$
- Kirchhoff potential law: there are $p(u)$s
- Ohm's law: $q_{(u,v)} = (p_u - p_v)/r_e$

Dynamics: $\dot{x}_e = |q_e| - x_e$

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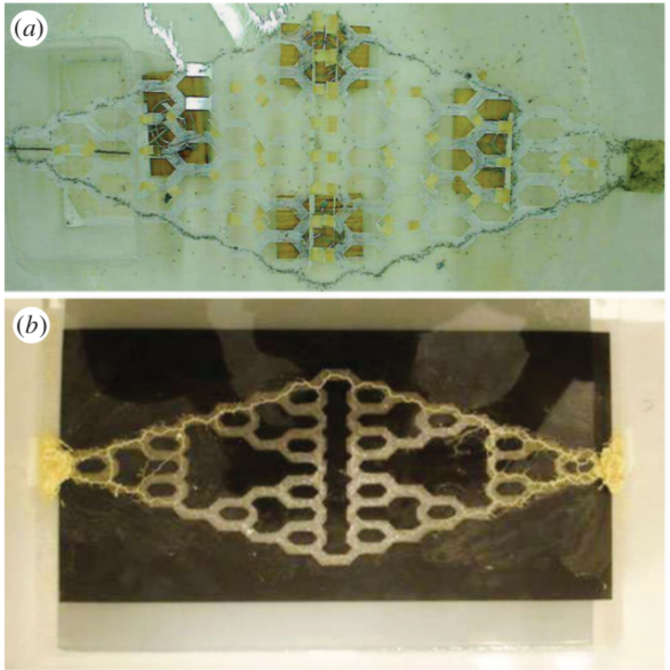
Ma, Johansson, Tero, Nakagaki and Sumpter, J. of the Royal Society Interface '13



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flow of ants from u to v = $\dfrac{\#\{\text{ants in } u\} - \#\{\text{ants in } v\}}{\text{lenght}/\text{pheromone}}$

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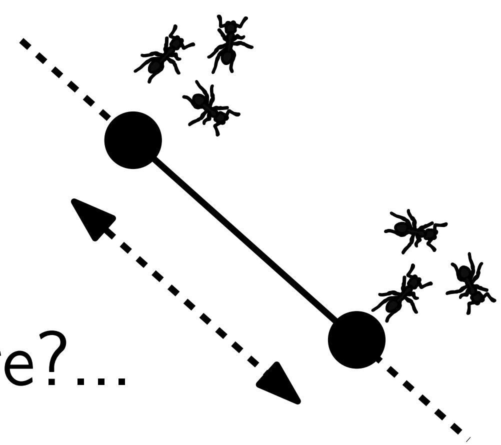

no ant traversing edge?...

Physarum Dynamics as an Algorithm

Bonifaci, Mehlhorn and Varma SODA'12:
Physarum dynamics converges on all graphs
(elegant proof in Bonifaci IPL'13)

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Euler's discretization

$$x(t+1) - x(t) = h(|q(t)| - x(t))$$

Becchetti, Bonifaci, Dirnberger, Karrenbauer and Mehlhorn ICALP'13:
Discretized physarum computes $(1+\epsilon)$-apx.
in $\mathcal{O}(mL(\log n + \log L)/\epsilon^3)$

More Research on Physarum

Many sequels in TCS

- Bonifaci IPL'13,
- Straszak and Vishnoi ITCS'16,
- Straszak and Vishnoi SODA'16
- Becker et al. ESA'17
- ...

More Research on Physarum

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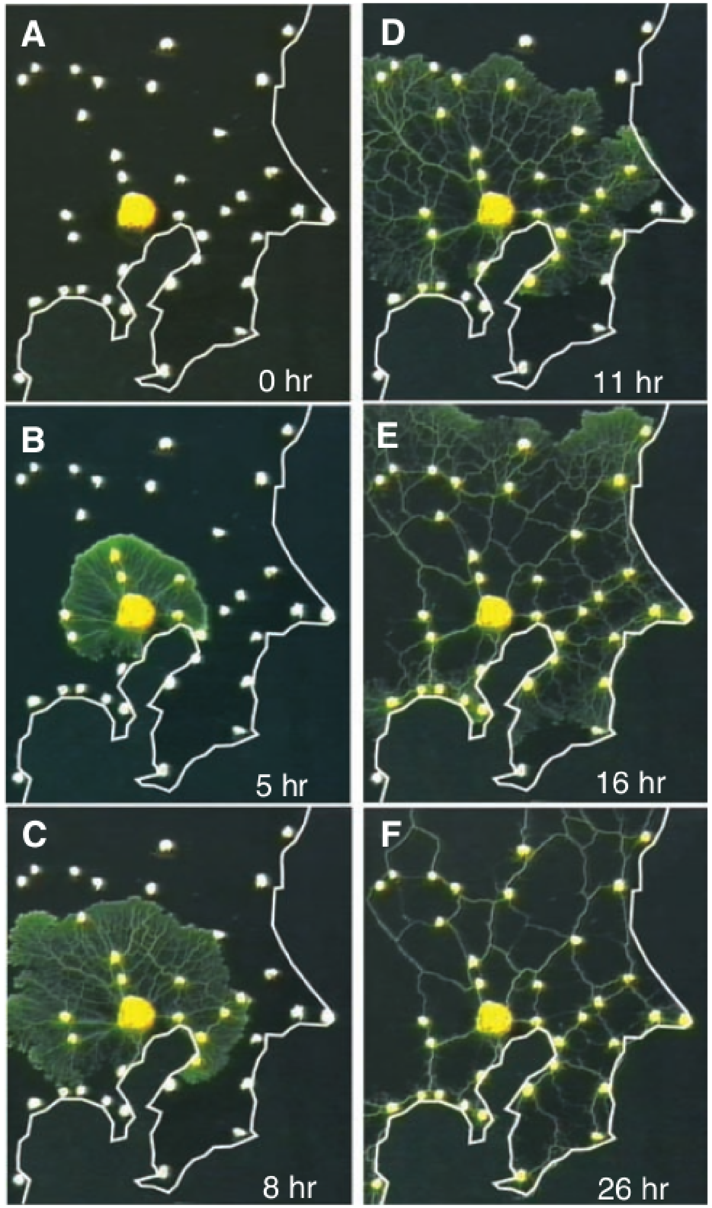
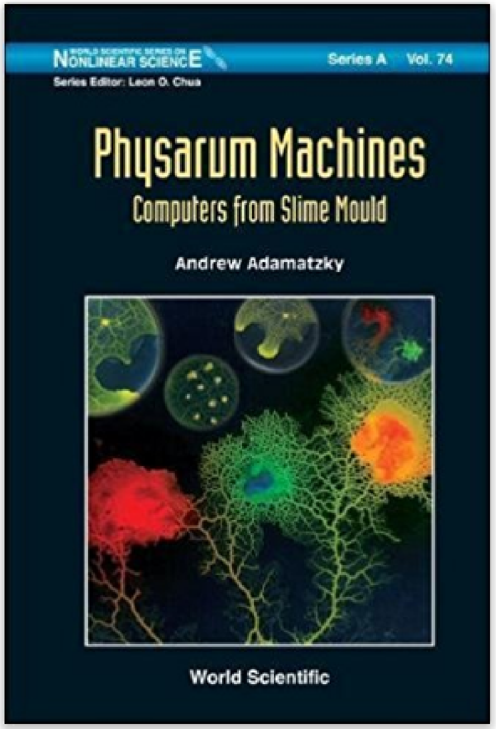
- ...

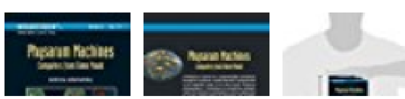
Some sequels elsewhere...
Tero et al. Science 2010: Physarum re-builds Tokyo's rail network!



More Research on Physarum

Look inside





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Physarum Machines: Computers from Slime Mould (World Scientifi Nonlinear Science, Series A) Hardcover – August 26, 2010

by Andrew Adamatzky (Author)

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Hardcover	Paperback
$60.82	$104.00
30 Used from $62.10	5 Used from $99.66
34 New from $60.82	4 New from $99.66

A Physarum machine is a programmable amorphous biological computer experimentally implemented in the vegetative state of true slime mould Physarum polycephalum. It comprises an amorphous yellowish mass with networks of protoplasmic veins, programmed by spatial configurations of attracting and repelling gradients.
This book demonstrates how to create experimental Physarum machines for computational geometry and optimization, distributed manipulation and transportation, and general-purpose computation. Being very cheap to make and easy to maintain, the machine also functions on a wide range of substrates and in a broad scope of environmental conditions. As such a Physarum machine is a green and environmentally friendly unconventional computer.

The book is readily accessible to a nonprofessional reader, and is a priceless source of experimental tips and inventive theoretical ideas for anyone who is inspired by novel and emerging non-silicon computers and robots.

Read less

How to Compute with Electrical Networks

Physarum have to solve Kirchhoff's equations

$\sum_{v \sim u} q_{(u,v)} = \sum_{v \sim u} (p_u - p_v)/r_e = b(u)$

or $\boxed{Lp = b}$

- edge's weight x_e/ℓ_e
- D diagonal matrix of nodes' volumes
- A weighted incidence matrix
- $L = D - A$

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Previous approaches: centralized computation

- Can be accomplished if every node is agent that follows elementary protocol?
(*biologically*: what happens microscopically?)
- If yes, what is convergence time and communication overhead?

Distributed Jacobi's Method

Jacobi's iterative method (Varga, 2009):
Bound on convergence rate w.r.t. *graph conductance*
exploiting structure of laplacian
(cfr. also DeGroot's model)

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$$Lp = (D - A)p = b \implies p = \underbrace{D^{-1}A}_{P}p + D^{-1}b$$

Jacobi's matrix = transition matrix

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Jacobi's: $\tilde{p}(t+1) = P\tilde{p}(t) + b$

Error $e(t) = p - \tilde{p}(t) = e_{\perp}(t) + \alpha \mathbf{1}$
(p *doesn't care about* α: $L\mathbf{1} = 0$!)

Jacobi's Convergence

The new error is

$$p - e_\perp(t+1) - \alpha(t+1) \cdot \mathbf{1} = \tilde{p}(t+1)$$

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$$= D^{-1}\left(A\tilde{p}(t) + b\right)$$

Jacobi's Convergence

The new error is

$$\begin{aligned} p - e_\perp(t+1) - \alpha(t+1)\cdot\mathbf{1} &= \tilde{p}(t+1) \\ &= D^{-1}\left(A\tilde{p}(t) + b\right) \\ &= D^{-1}\left(A\left(p - e_\perp(t) - \alpha(t)\cdot\mathbf{1}\right) + b\right) \end{aligned}$$

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$$
\begin{aligned}
p - e_{\perp}(t+1) - \alpha(t+1) \cdot \mathbf{1} = \tilde{p}(t+1) \\
&= D^{-1}\left(A\tilde{p}(t) + b\right) \\
&= D^{-1}\left(A\left(p - e_{\perp}(t) - \alpha(t) \cdot \mathbf{1}\right) + b\right) \\
&= p - Pe_{\perp}(t) - \alpha(t) \cdot \mathbf{1},
\end{aligned}
$$

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thus $e_{\perp}(t+1) = Pe_{\perp}(t) - (\alpha(t+1) - \alpha(t))\mathbf{1}$

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thus $e_{\perp}(t+1) = Pe_{\perp}(t) - (\alpha(t+1) - \alpha(t))\mathbf{1}$

and $e_{\perp}(t) = \left(I - \frac{1}{n}\mathbf{1}\mathbf{1}^{\mathsf{T}}\right)P^{t}e_{\perp}(0)$.

Jacobi's Convergence

$P = D^{-1}A$ is similar to $N = D^{-1/2}AD^{-1/2}$.
Thus

$$P^t = (D^{-1}A)^t = (D^{-\frac{1}{2}}ND^{\frac{1}{2}})^t = D^{-\frac{1}{2}}N^tD^{\frac{1}{2}}.$$

Observe that

- N has n orthonormal eigenvec. $\vec{x}_1, \ldots, \vec{x}_n$, corresponding to eigenvectors $\vec{y}_1, \ldots, \vec{y}_n$ of P via $\vec{x}_i = D^{1/2}\vec{y}_i$ for each i.
- Both $\vec{x}_i$ and $\vec{y}_i$, for each i, are associated to the same eigenvalue ρ_i of P.

Jacobi's Convergence

$$\|e_\perp(t)\|$$
$$= \left\|\left(I - \tfrac{1}{n}\mathbf{1}\mathbf{1}^\intercal\right) P^t e_\perp(0)\right\|$$
$$= \left\|\left(I - \tfrac{1}{n}\mathbf{1}\mathbf{1}^\intercal\right) D^{-\frac{1}{2}} N^t D^{\frac{1}{2}} e_\perp(0)\right\|$$

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&= \left\|\left(I - \tfrac{1}{n}\mathbf{1}\mathbf{1}^\intercal\right) D^{-\frac{1}{2}} \left(\textstyle\sum_{i=2}^{n} \rho_i^t \vec{x}_i \vec{x}_i^\intercal\right) D^{\frac{1}{2}} e_\perp(0)\right\|
\end{aligned}
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&\leq \left\|\left(I - \tfrac{1}{n}\mathbf{1}\mathbf{1}^\intercal\right)\right\| \left\|D^{-\frac{1}{2}}\right\| \left\|\textstyle\sum_{i=2}^n \rho_i^t \vec{x}_i \vec{x}_i^\intercal\right\| \left\|D^{\frac{1}{2}}\right\| \|e_\perp(0)\| \\
&\leq \sqrt{\tfrac{\mathrm{vol}_{\max}}{\mathrm{vol}_{\min}}}\, \max(|\rho_2|, |\rho_n|)^t \|e_\perp(0)\|,
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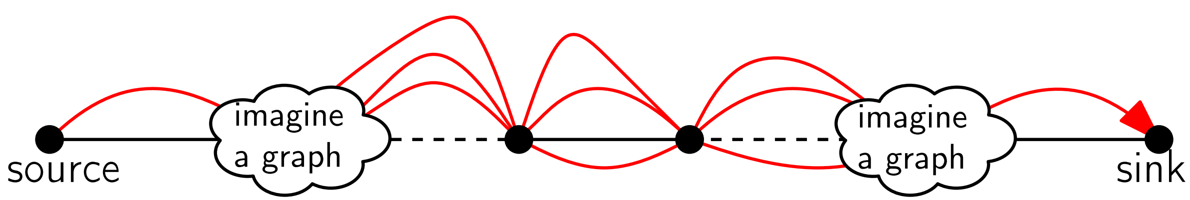
↑ Conductance by Cheeger's inequality

Randomized Token Diffusion Process

Doyle and Snell, ’84 & Tetali, ’91:

Times a random walk transits through given edge *until hitting* the sink

- *global* requirement
- no accuracy and msg. complexity bounds



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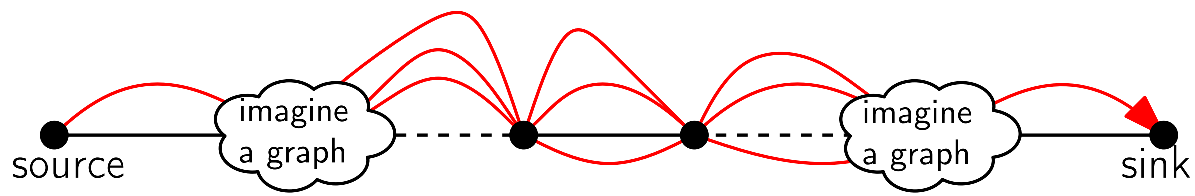

Our’s:

How many tokens are on a node

- *local* requirement
- accuracy and msg. complexity w.r.t. *edge expansion*

Randomized Token Diffusion Process

Process

- At the beginning of each step, K new tokens *appear* at the source
- Each token independently performs a *weighted random walk* at each step
- Each token that hits the sink *disappears*

Estimator

$V_K^{(t)} = \frac{Z_K^{(t)}(u)}{K \cdot \mathrm{vol}(u)}$ where $Z_K^{(t)}(u)$ number of tokens on u

Expected Behavior

Define inductively $\mathbf{p}^{(t)}$ by

$$p_u^{(0)} = 0, \qquad \text{for all } u \in \mathcal{V},$$

$$p_u^{(t+1)} = \begin{cases} \frac{1}{\mathrm{vol}(u)} \left(\sum_{v \sim u} w_{uv} p_v^{(t)} + b_u \right) & \text{if } u \neq \mathrm{sink}, \\ 0 & \text{if } u = \mathrm{sink}. \end{cases}$$

Lemma. If $V_K^{(t)}(u) = \frac{Z_K^{t)}(u)}{K\mathrm{vol}(u)}$, then $\mathbb{E}[V_K^{(t)}(u)] = p_u^{(t)}$.

Correctness of Token Diffusion

We can write

$$\begin{cases} \mathbf{p}^{(0)} & = \vec{0}, \\ \mathbf{p}^{(t+1)} & = \underline{P}\,\mathbf{p}^{(t)} + D^{-1}\underline{\vec{b}}, \end{cases}$$

with $\underline{P}$ and $\underline{\vec{b}}$ obtained by zeroing out entries on row and column of sink.

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Lemma. The spectral radius of $\underline{P}$, $\underline{\rho}$, satisfies $\underline{\rho} = 1 - \sum_{i=1}^{n} v_i \cdot P_{i,\text{sink}} / \|v_1\|$, where $\vec{v}_1$ is left Perron eigenvector of $\underline{P}$.

Correctness of Token Diffusion

We can write

$$\begin{cases} \mathbf{p}^{(0)} & = \vec{0}, \\ \mathbf{p}^{(t+1)} & = \underline{P}\,\mathbf{p}^{(t)} + D^{-1}\underline{\vec{b}}, \end{cases}$$

with $\underline{P}$ and $\underline{\vec{b}}$ obtained by zeroing out entries on row and column of sink.

Lemma. The spectral radius of $\underline{P}$, $\underline{\rho}$, satisfies $\underline{\rho} = 1 - \sum_{i=1}^{n} v_i \cdot P_{i,\mathrm{sink}} / \|v_1\|$, where $\vec{v}_1$ is left Perron eigenvector of $\underline{P}$.

Theorem. System above converges to a valid potential with rate $\underline{\rho}$.

Time and Message Complexity

Theorem. $1 - \underline{\rho} \geq \frac{\overline{\lambda}_2}{2\text{vol}_{\max}(n-1)} \sum_i \frac{w_{in}}{w_{in} + \overline{\lambda}_2}$
where $\overline{\lambda}_2$ is 2nd smallest eigenvalue of non-normalized laplacian of graph with sink removed.

Connecting with *edge expansion*:
it is known
$\lambda_2(\mathcal{G}) \geq \text{vol}_{\max} - (\text{vol}_{\max}^2 - \theta(\mathcal{G})^2)^{1/2}$.

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Remark. As $t \to \infty$, the *expected message complexity* per round of Token Diffusion Algorithm is $O(K\, n\, \mathrm{vol}_{\max} \cdot E)$, where $E = \vec{p}^{\intercal} L \vec{p}$ is the *energy* of the electrical flow.

Stochastic Accuracy

X gives (ϵ, δ)-approximation of Y if
$\mathbf{P}\left(|X - Y| > \epsilon Y\right) \leq \delta$.

Lemma. For any K, $0 < \epsilon, \delta < 1$, t and u, such that $p_u^{(t)} \geq \frac{3}{\epsilon^2 K \mathrm{vol}(u)} \ln \frac{2}{\delta}$, the estimator provides an (ϵ, δ)-approximation of $p_u^{(t)}$.

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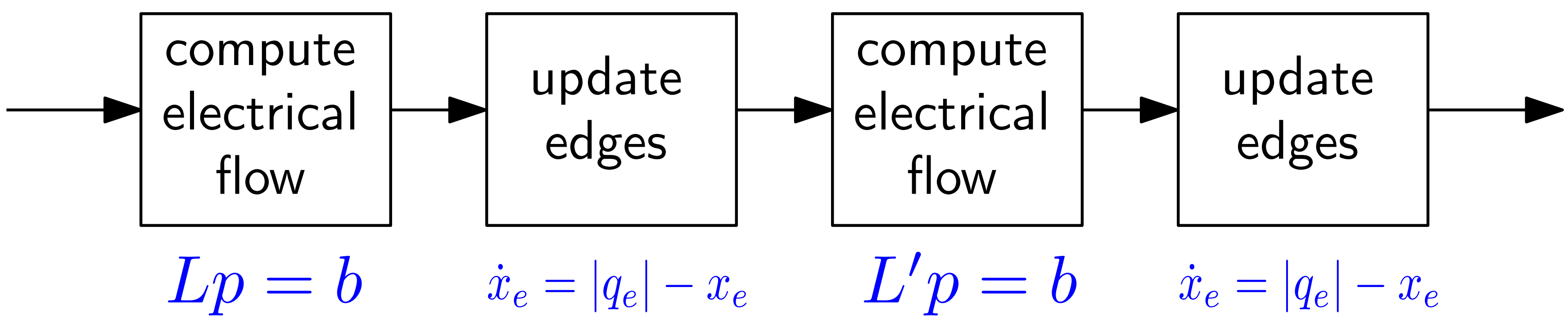
Proof. Chernoff bound requires $Y > 1/\epsilon^2$.

Open Problem: Analysis of *Distributed* Physarum?

Physarum dynamics et sim.:
Compute electrical flow, then *update edge-weigths*



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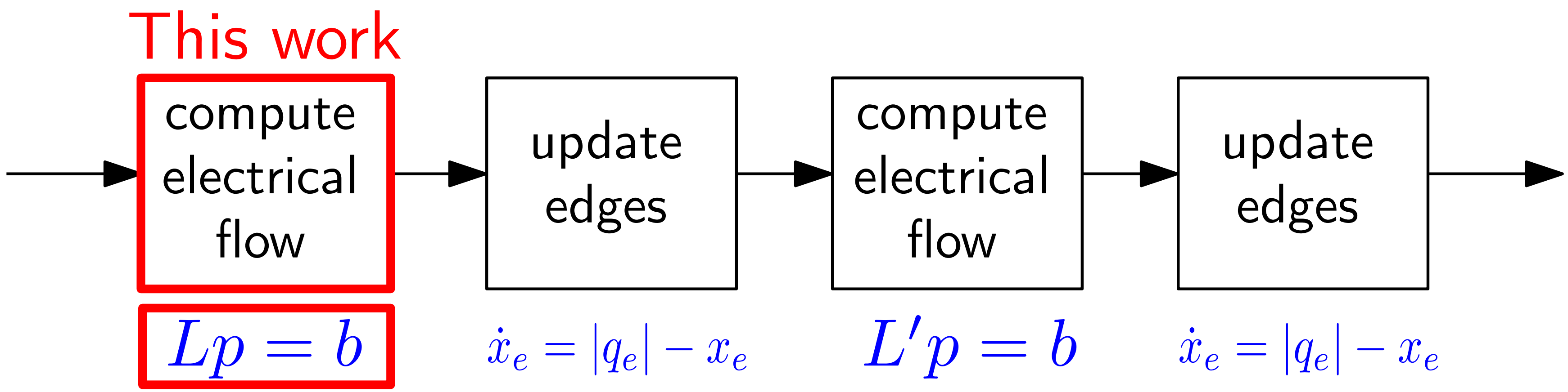

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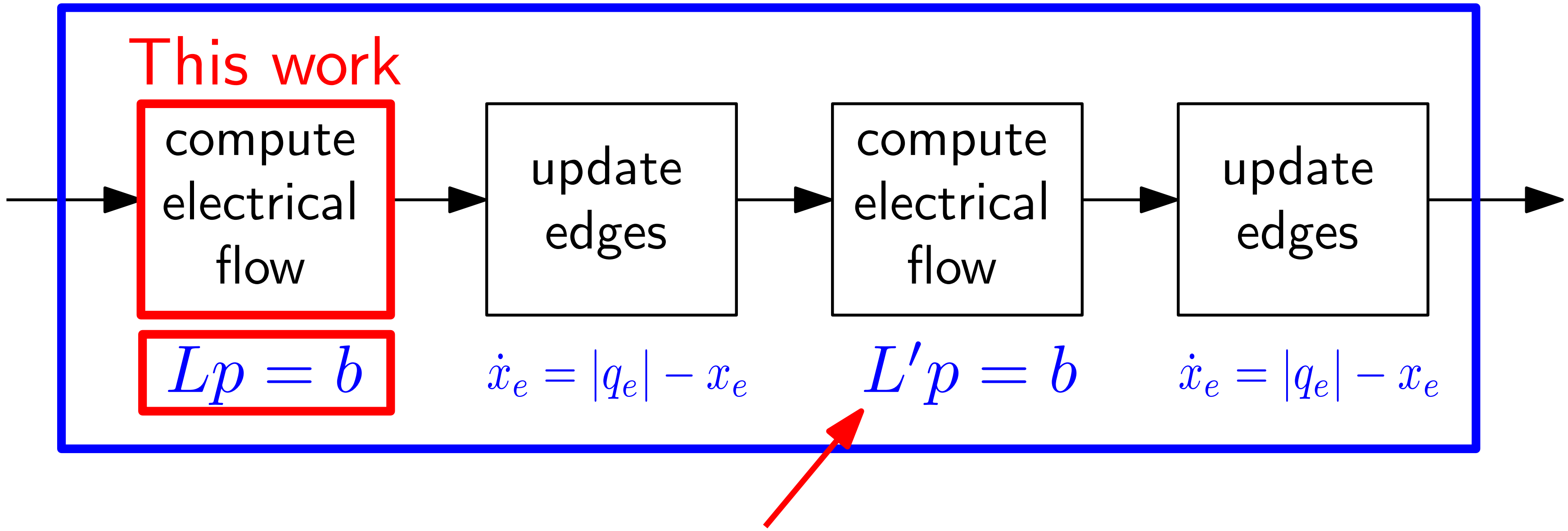

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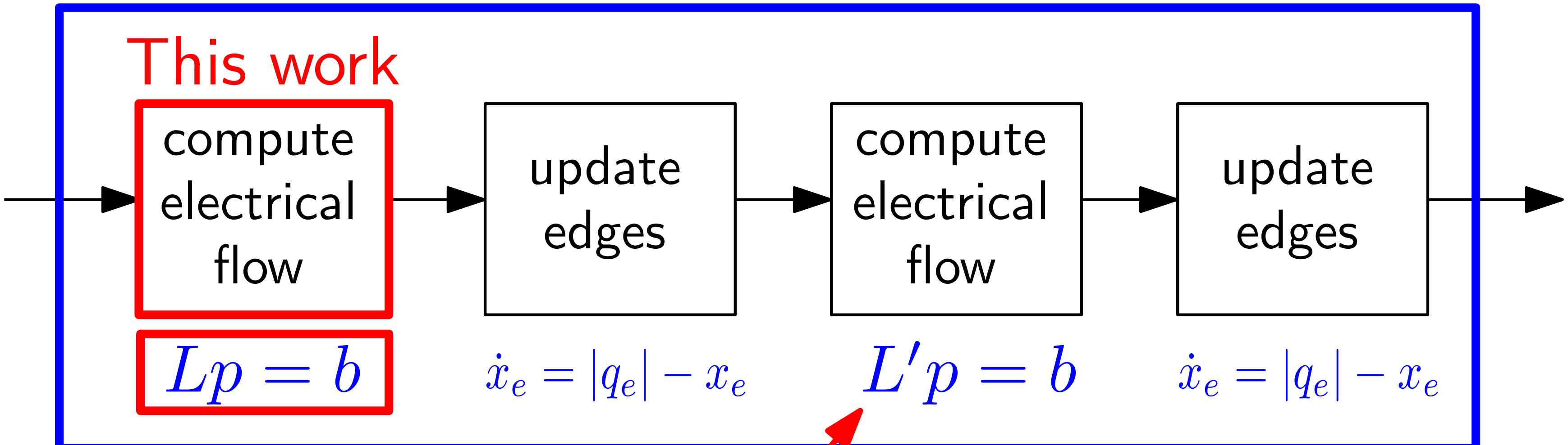

- Spectral structure of L'?
- Global convergence time?

Thank You!