

Self-Stabilizing Broadcast with $O(1)$ -Bit Messages^{*}

Emanuele Natale[†]

joint work with

Lucas Boczkowski^{*} and Amos Korman^{*}



SAPIENZA
UNIVERSITÀ DI ROMA

4th Workshop on Biological Distributed Algorithms
(BDA)

July 25-29, 2016

Chicago, Illinois

^{*}preprint at goo.gl/ETNc64

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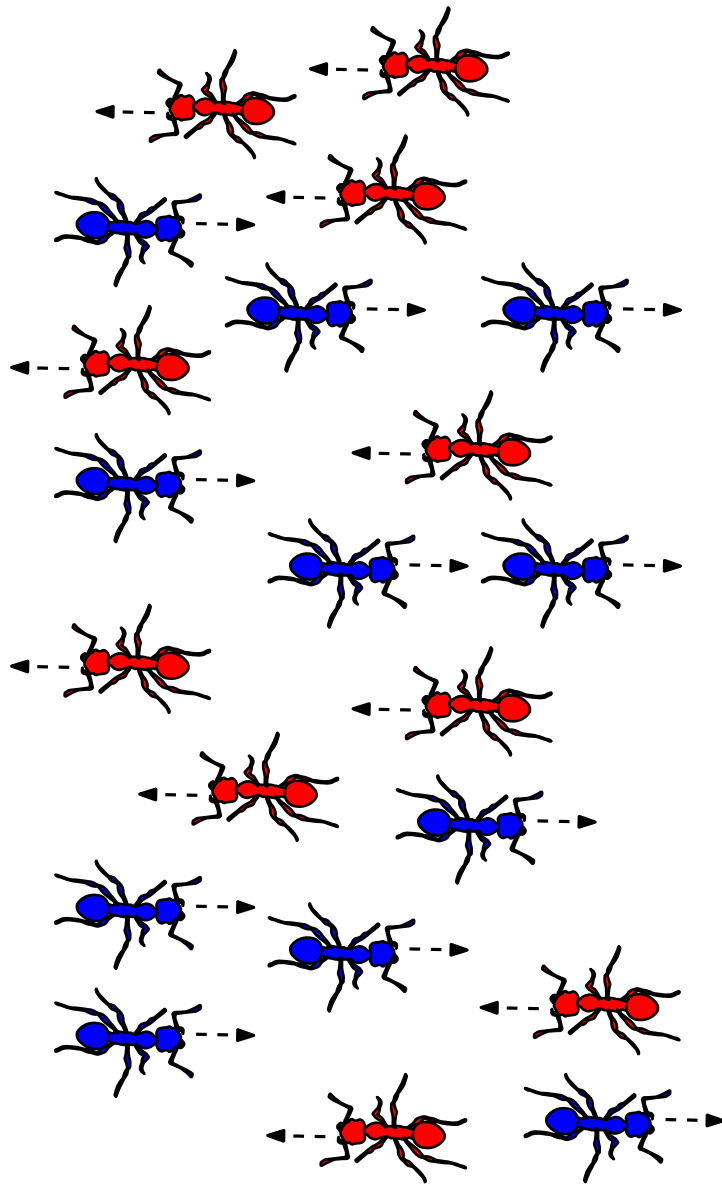
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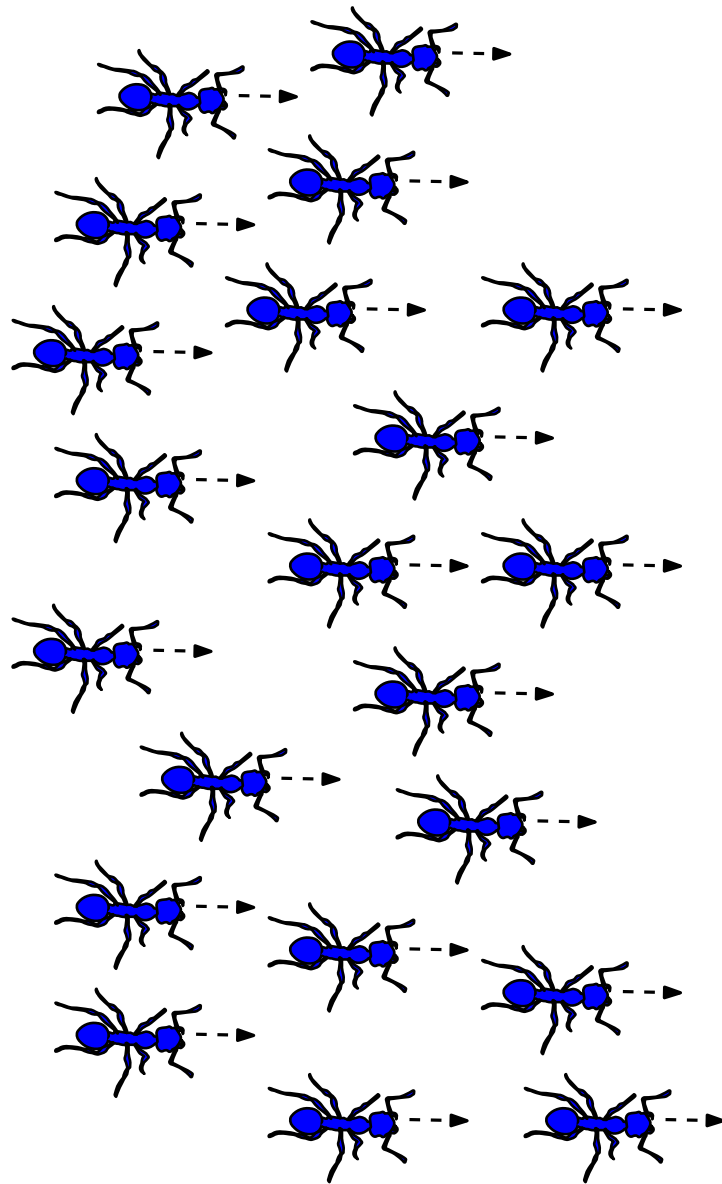
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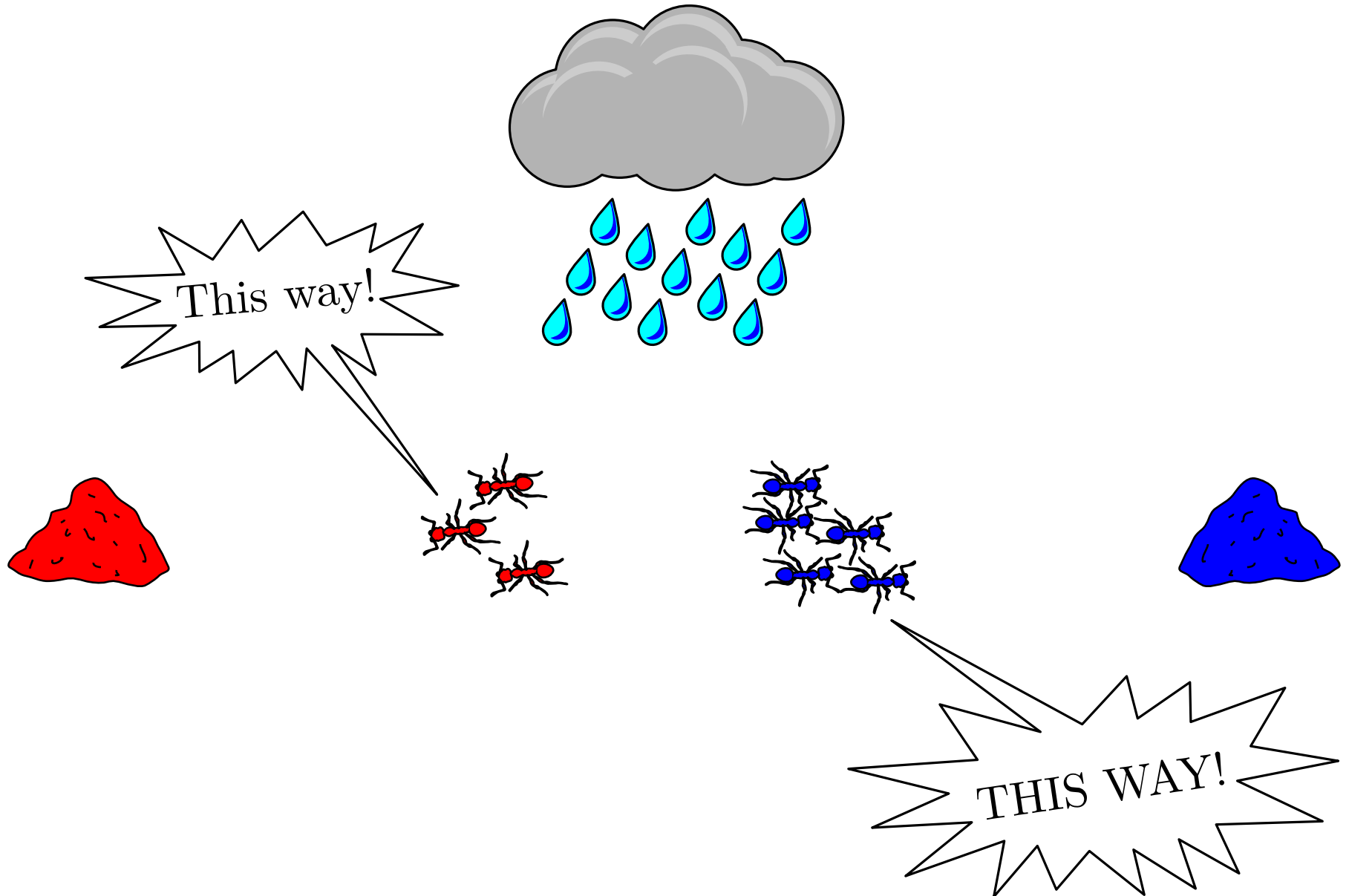
Bit Dissemination Problem



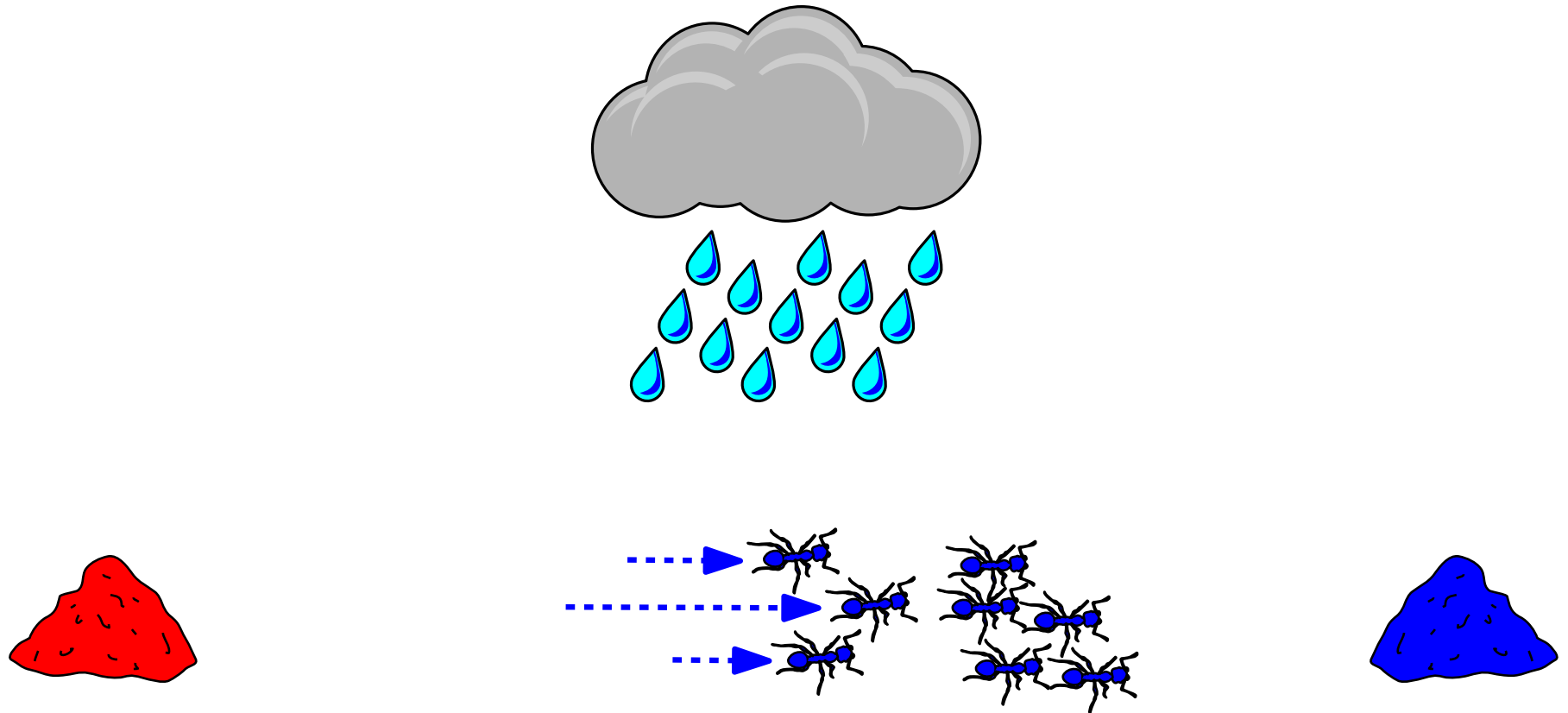
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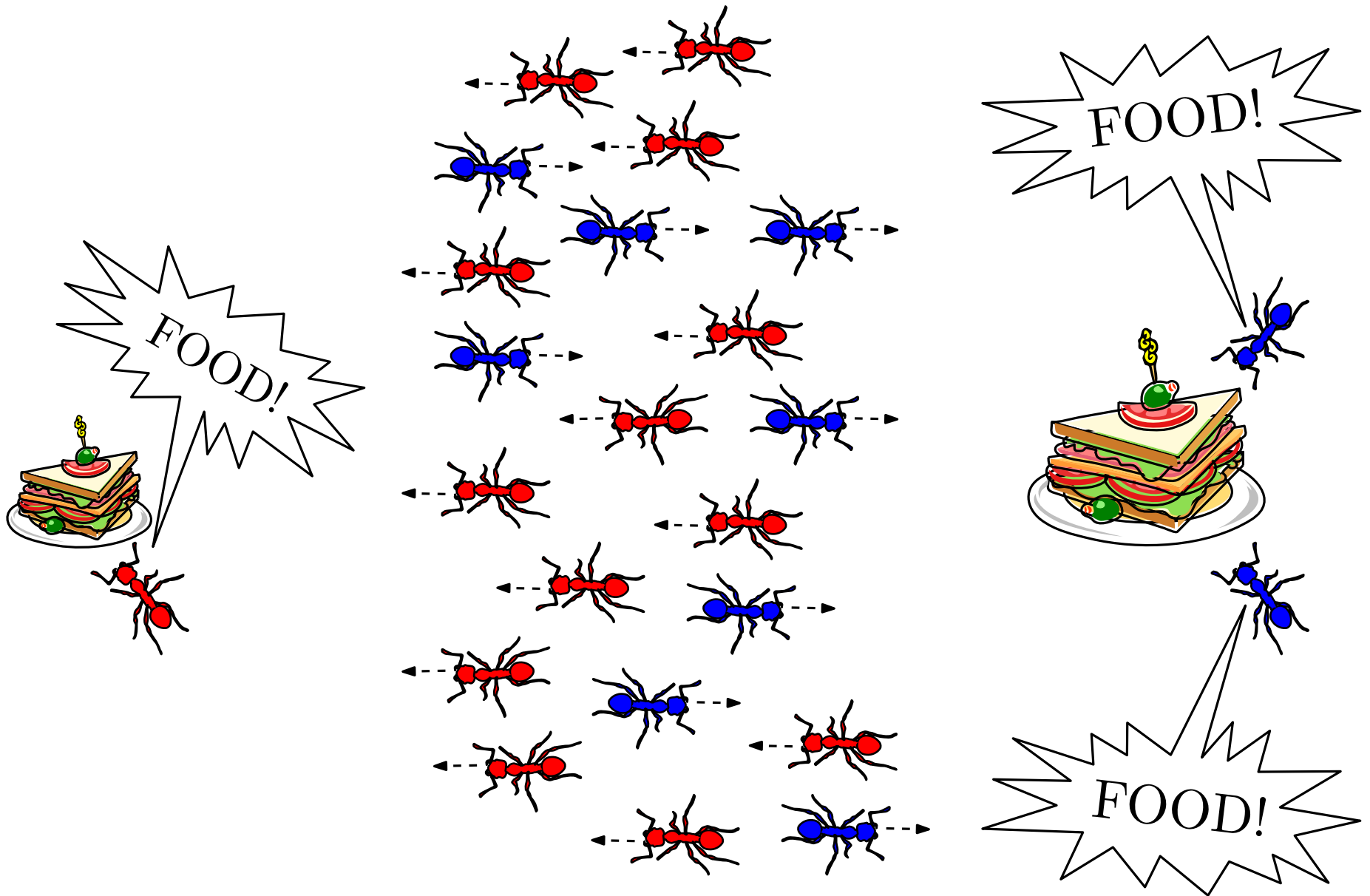
Majority Consensus Problem



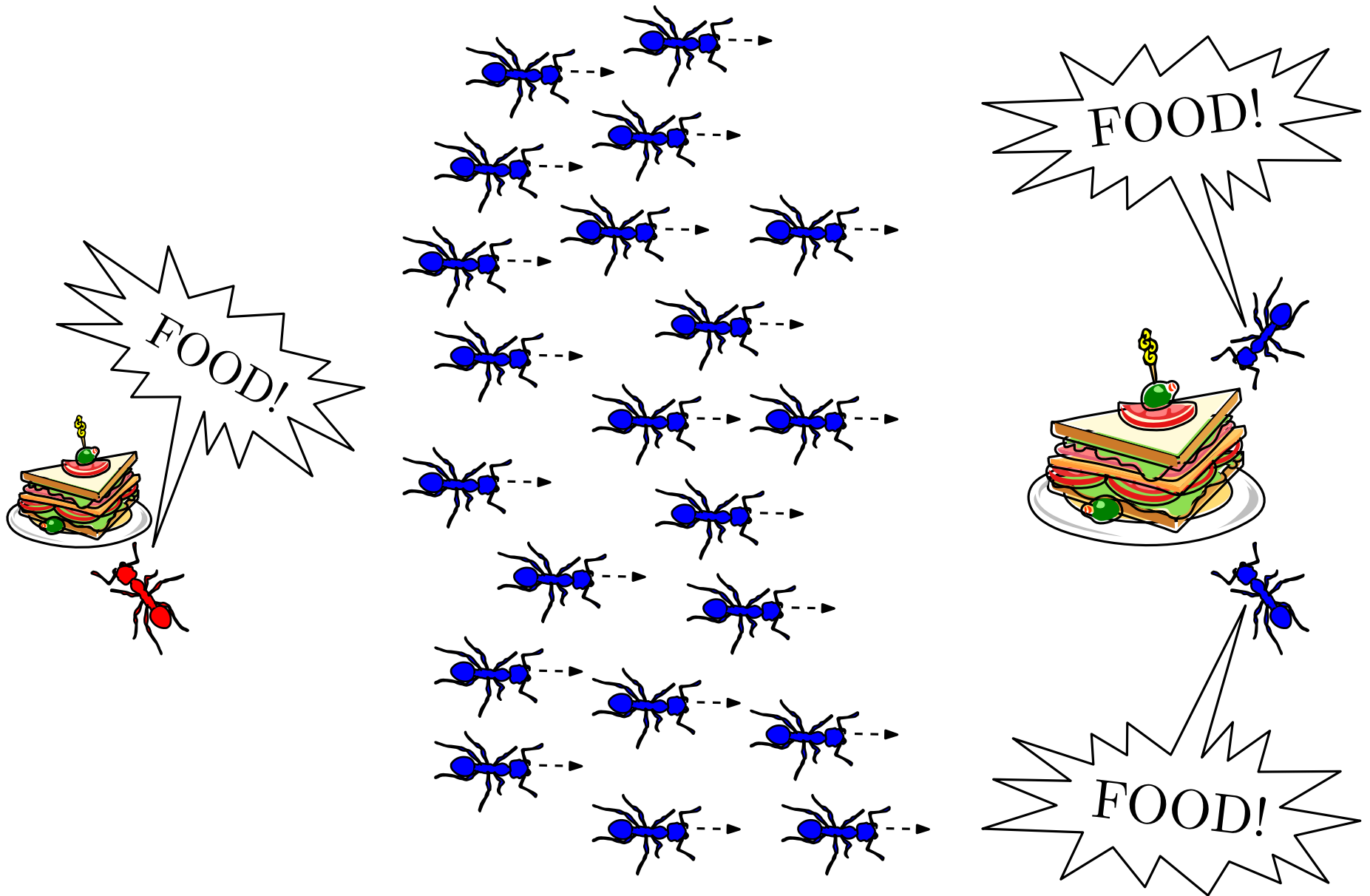
Majority Consensus Problem



Majority Bit Dissemination



Majority Bit Dissemination



Examples

Flocks of birds
[Ben-Shahar et al. '10]



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Schools of fish
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Examples

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Schools of fish
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Insects colonies
[Franks et al. '02]



Communication Model

Animal communication:

- Chaotic
- Anonymous
- Parsimonious
- Uni-directional (Passive/Active)
- Noisy

Communication Model

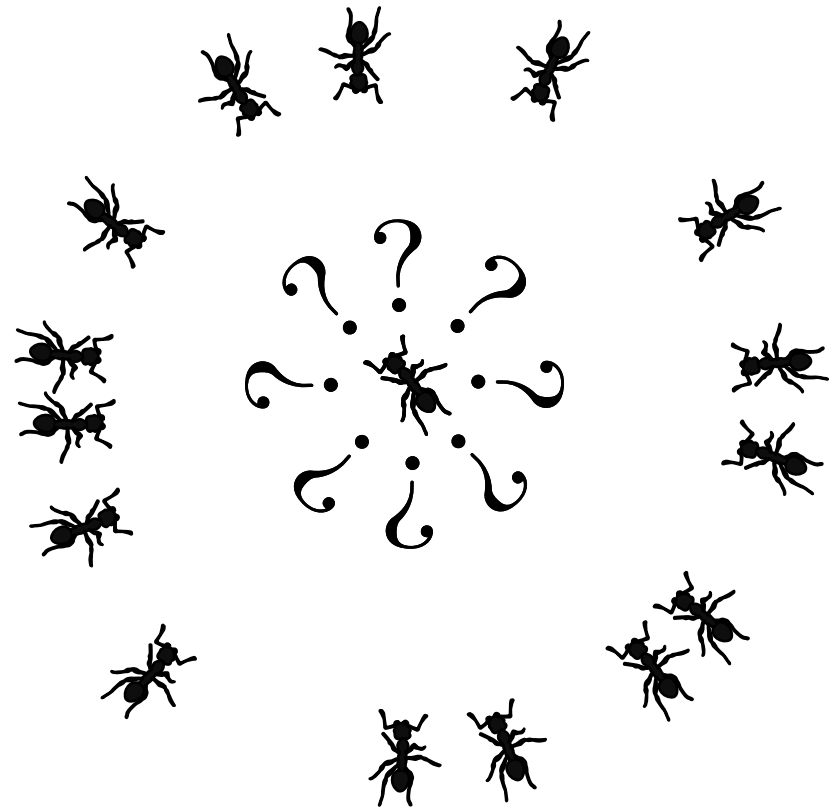
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$PUSH(h, \ell)$ model

[Demers '88]: at each round each agent can *send a ℓ -bit message* to h other agents chosen independently and uniformly at random.



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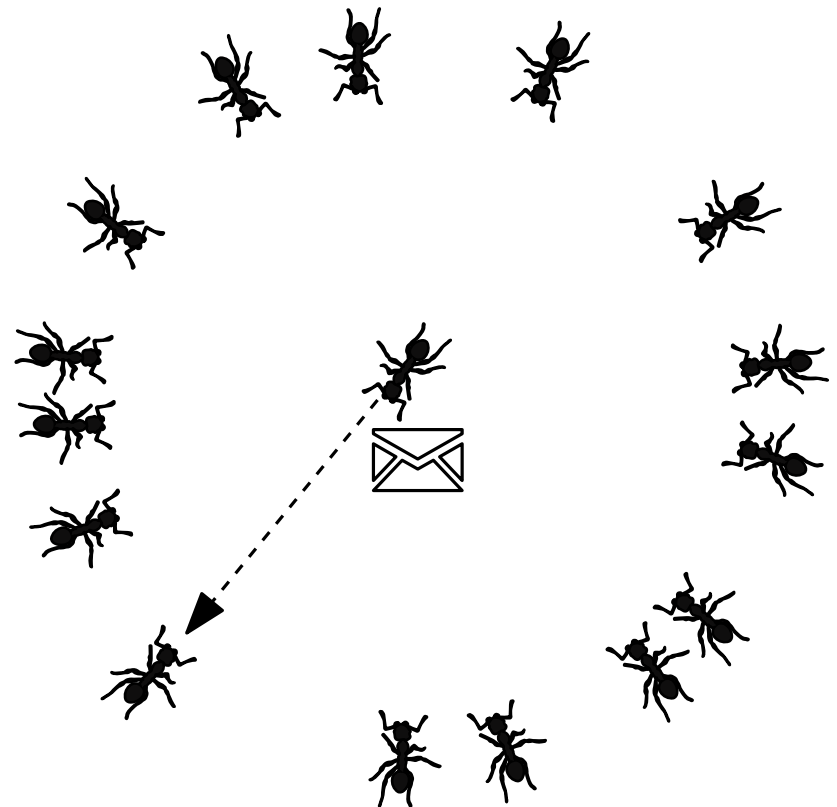
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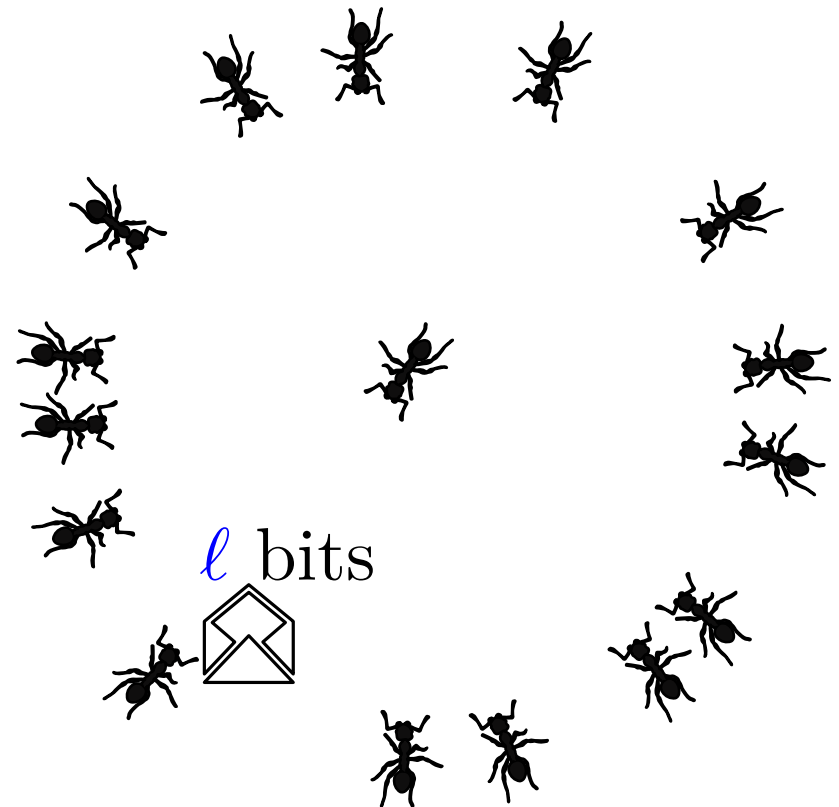
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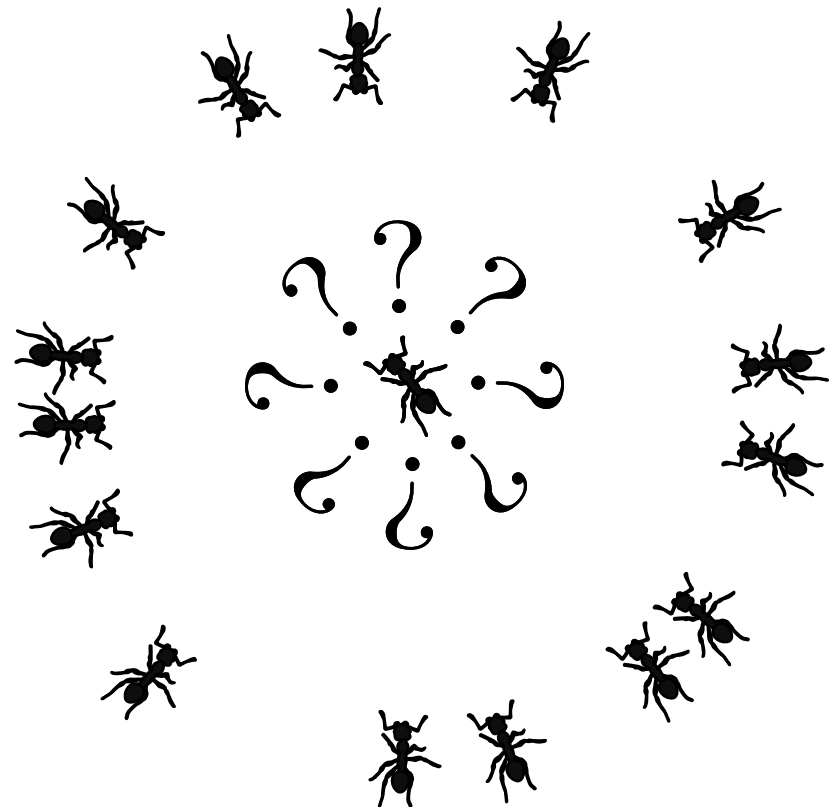
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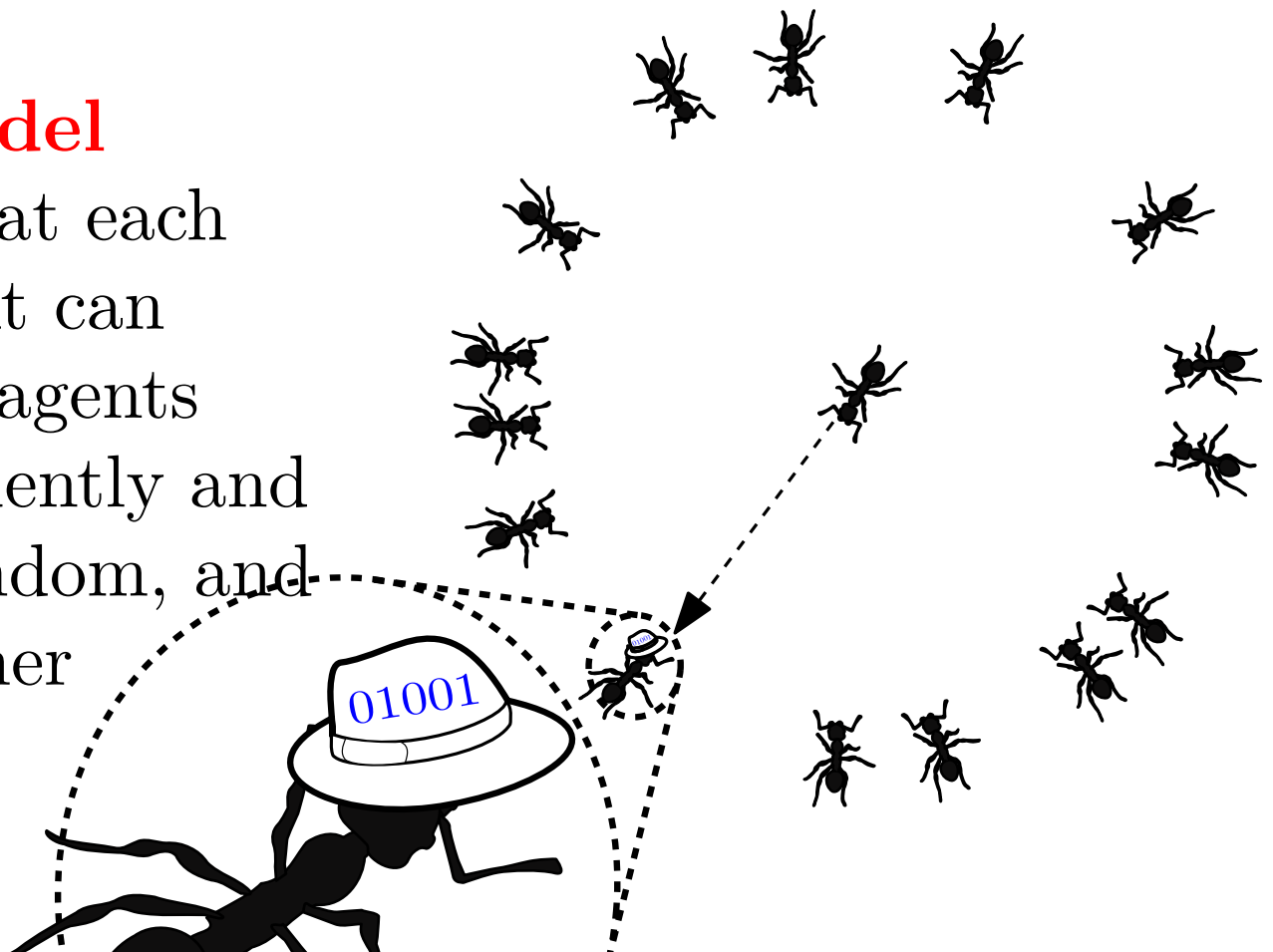
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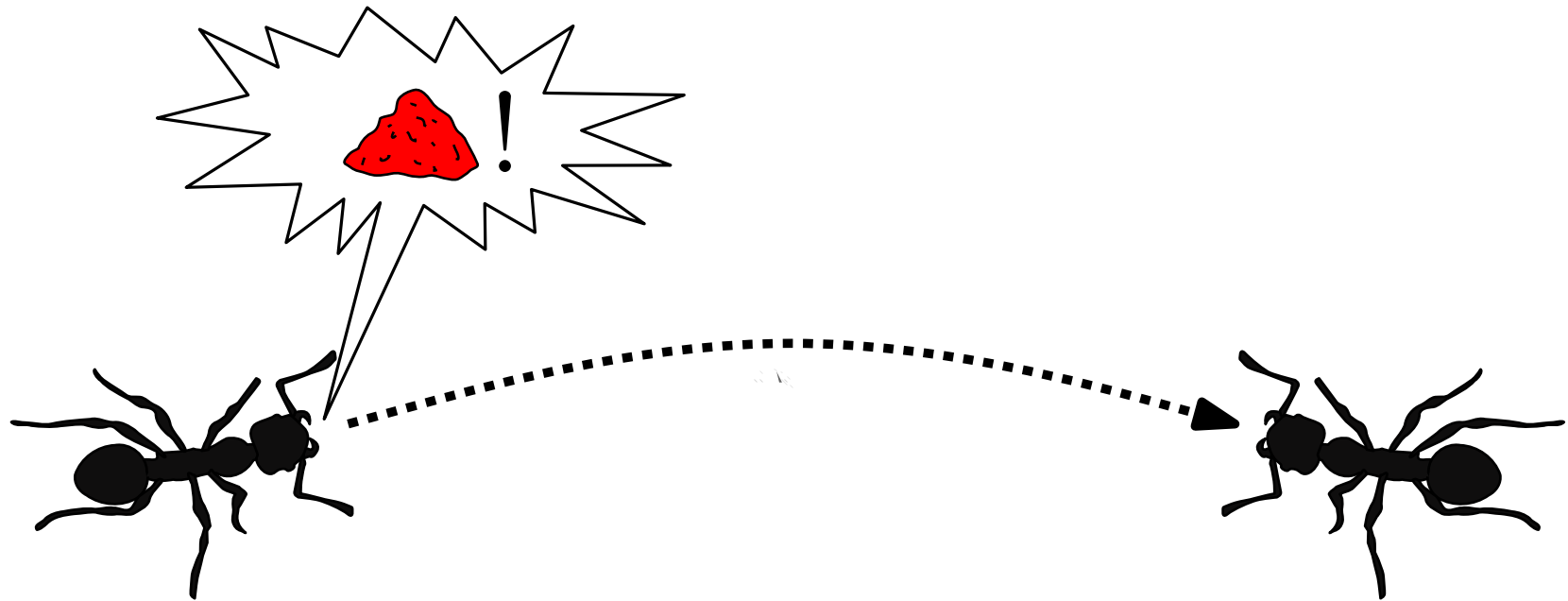
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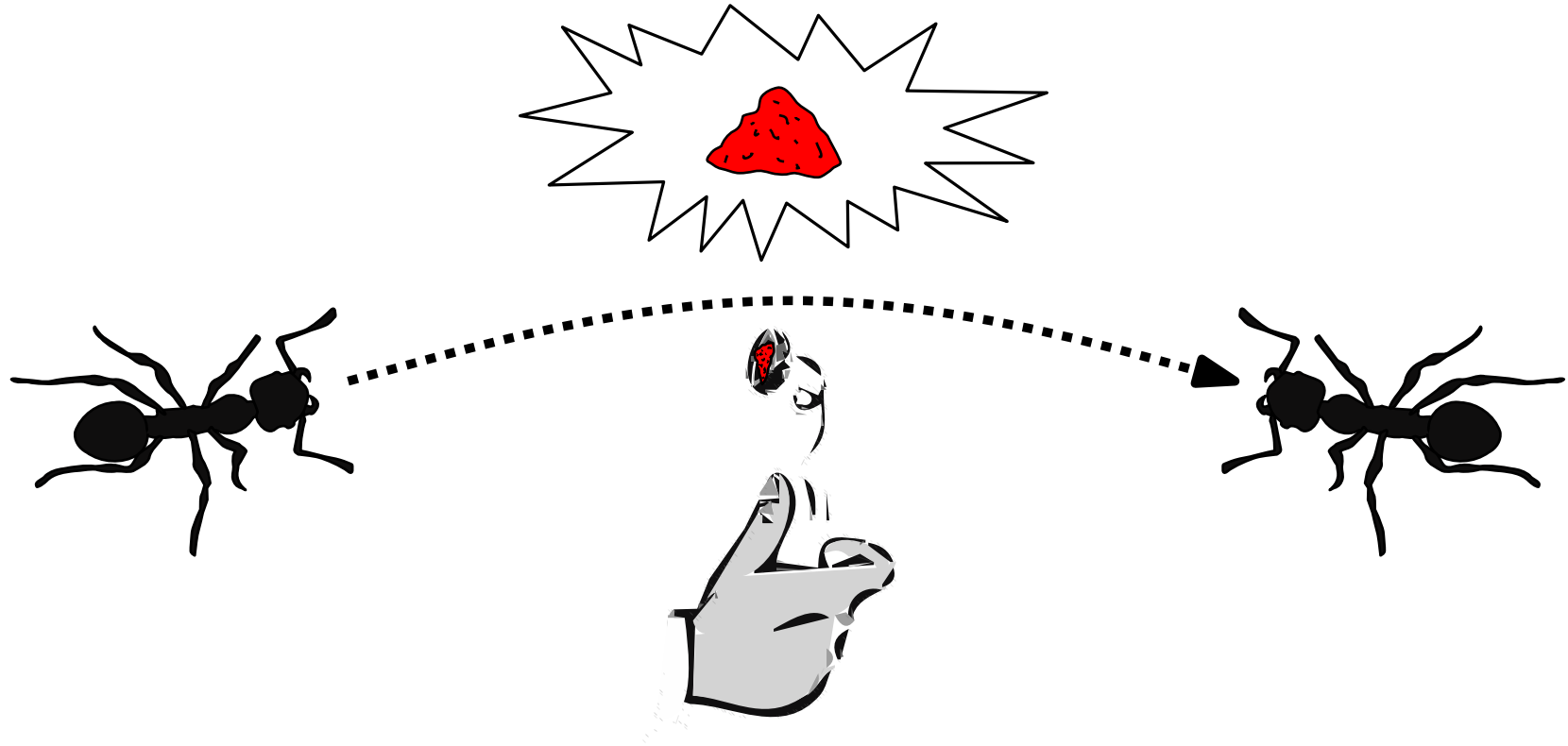
Related Work: MBD despite Noise

PUSH Model with *noise*: before being received, each bit is *flipped* with probability $1/2 - \epsilon$.



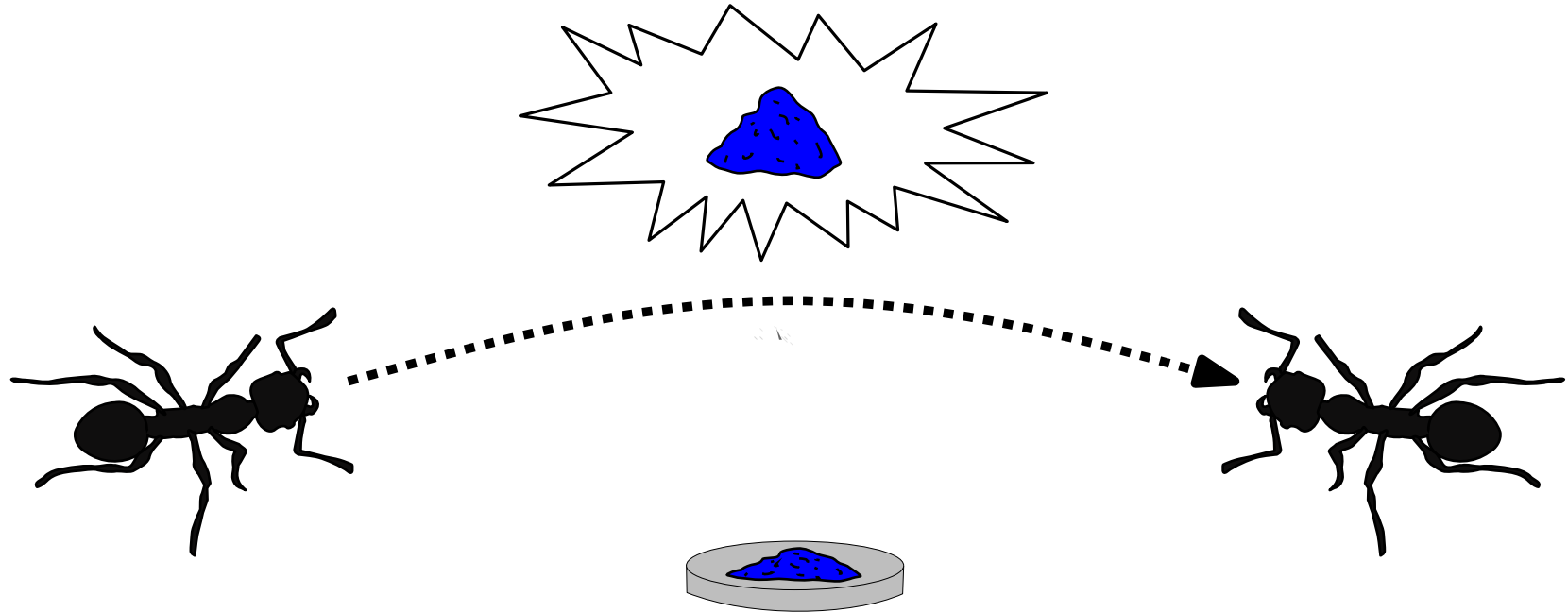
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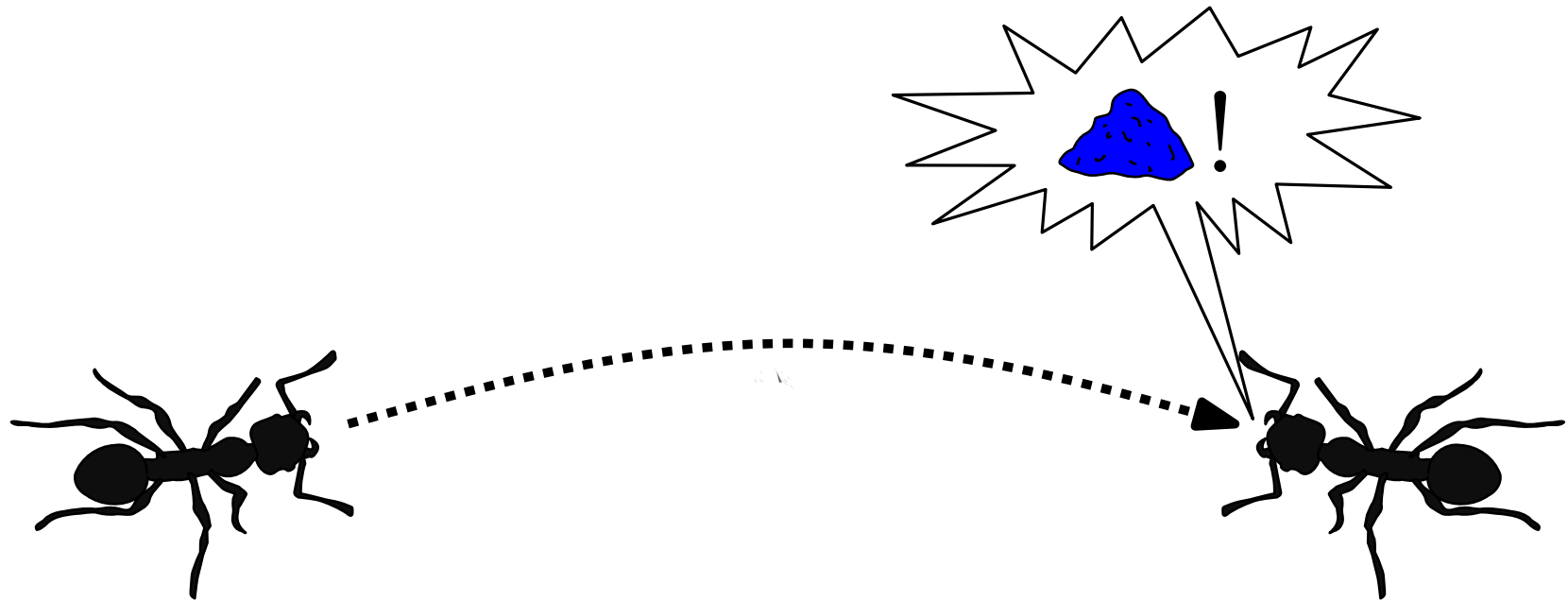
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O. Feinerman, B. Haeupler and A. Korman.

Breathe before speaking: efficient information dissemination despite noisy, limited and anonymous communication. (PODC '14)

⇒ Simple rules efficiently solve *binary* Majority Bit Dissemination despite noise.

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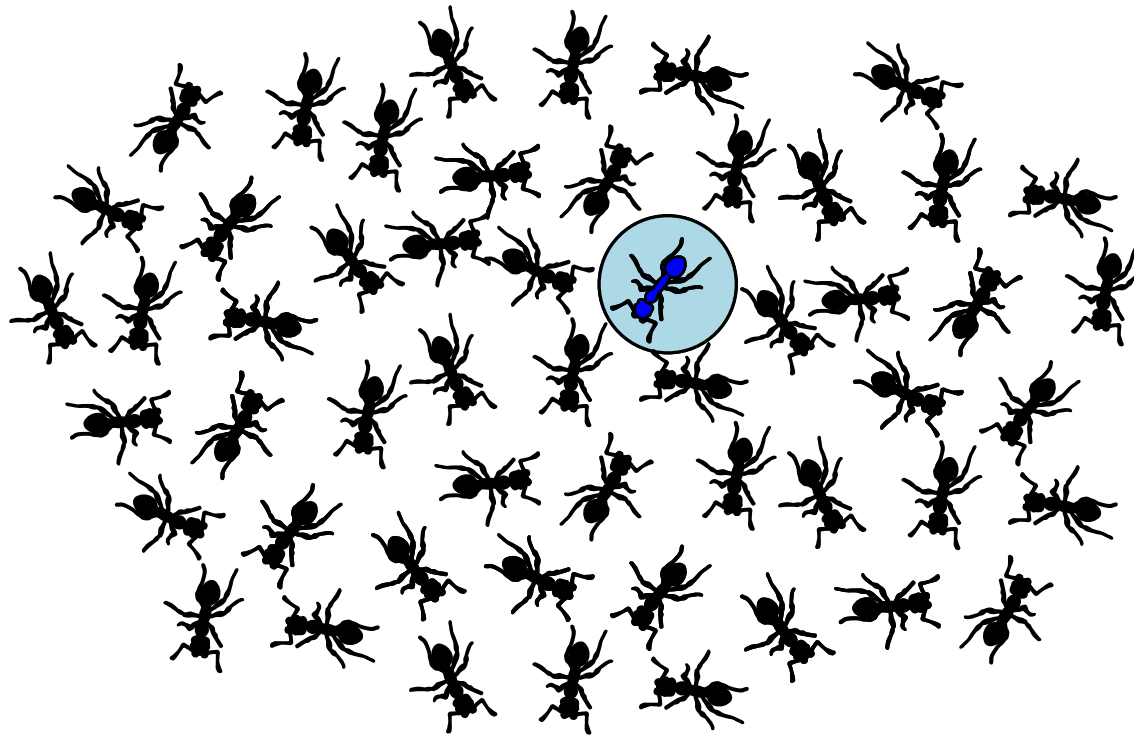
P. Fraigniaud, E. Natale.

Noisy Rumor-Spreading and Plurality Consensus.
(*BDA '15*, PODC '16)

⇒ Simple rules efficiently solve *multivalued* Plurality Opinion Dissemination despite noise.

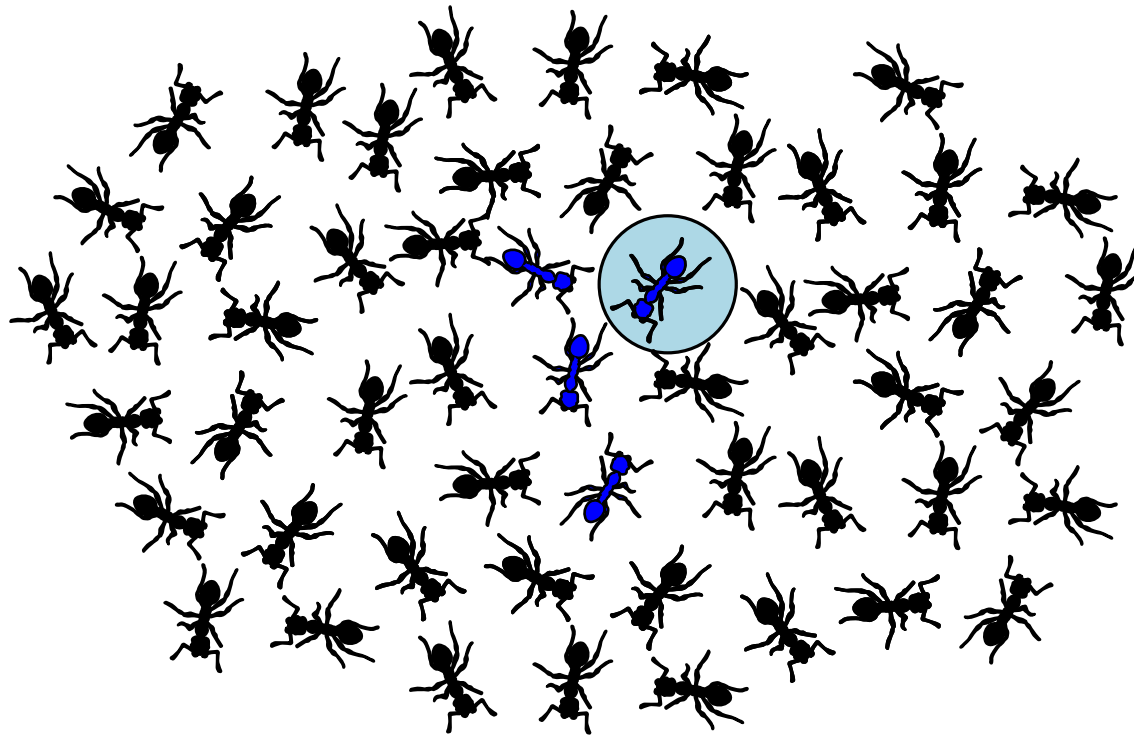
(Probabilistic) Self-Stabilization

Sources' bits (and other agents' states) may change in response to *external environment*



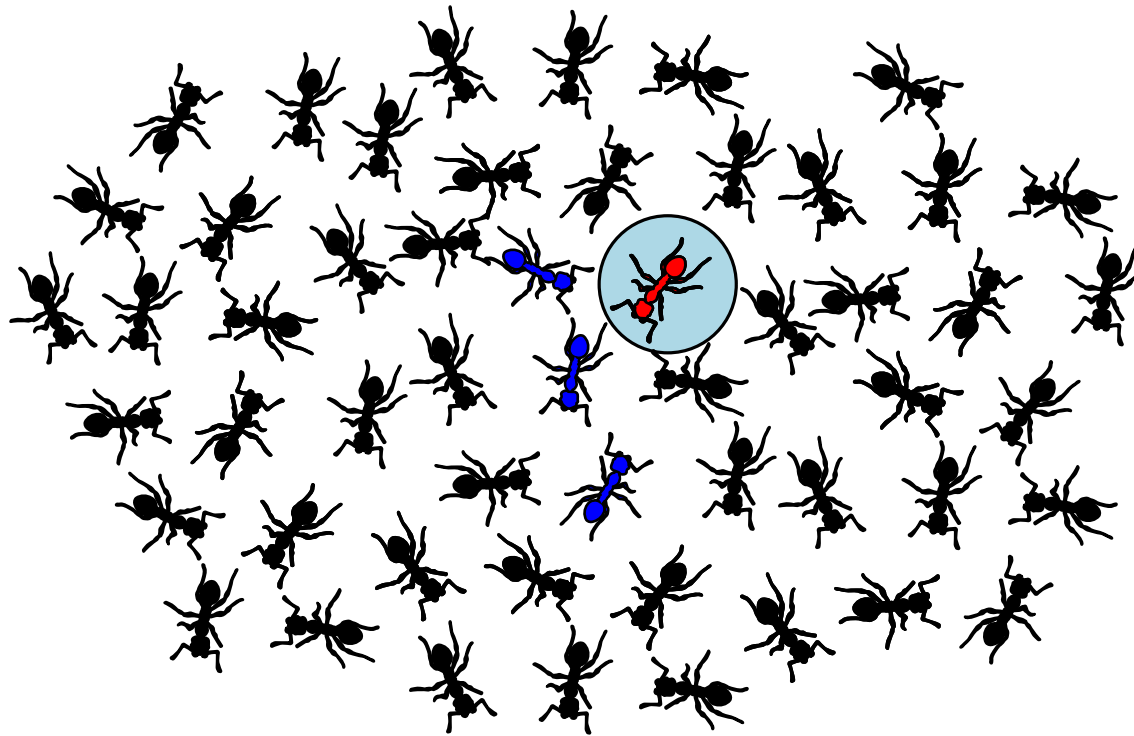
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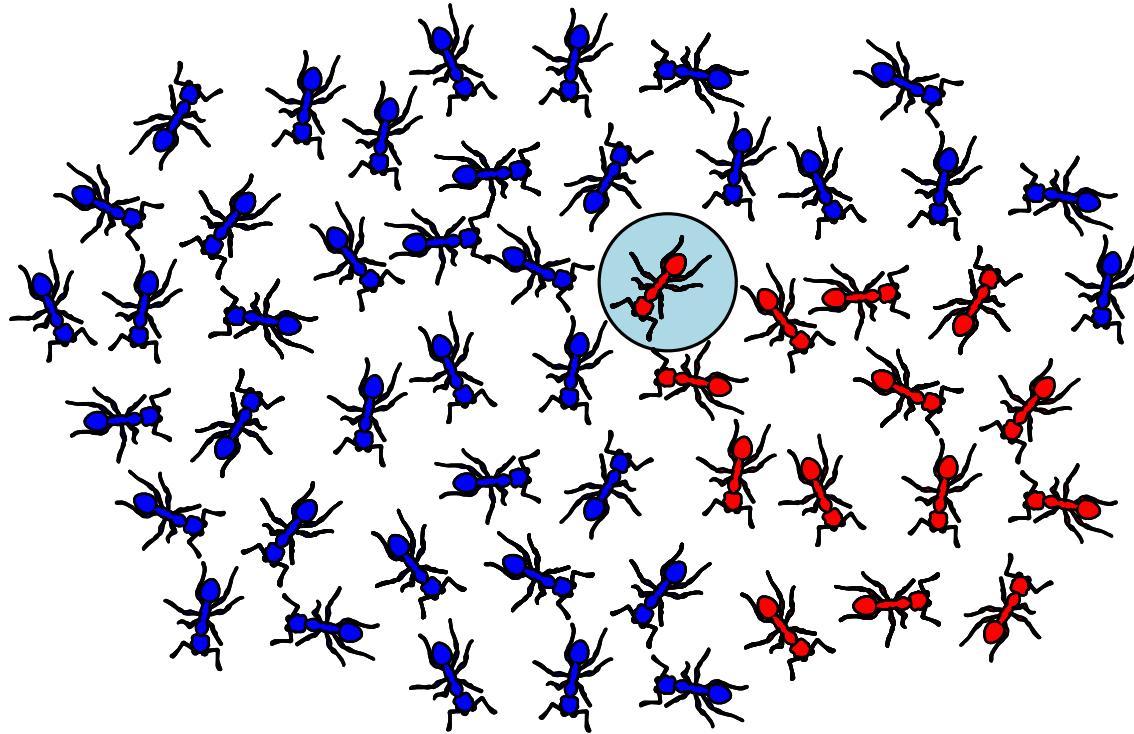
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blue vs red:

$39/14 \approx 2.8$



(Probabilistic) Self-Stabilization

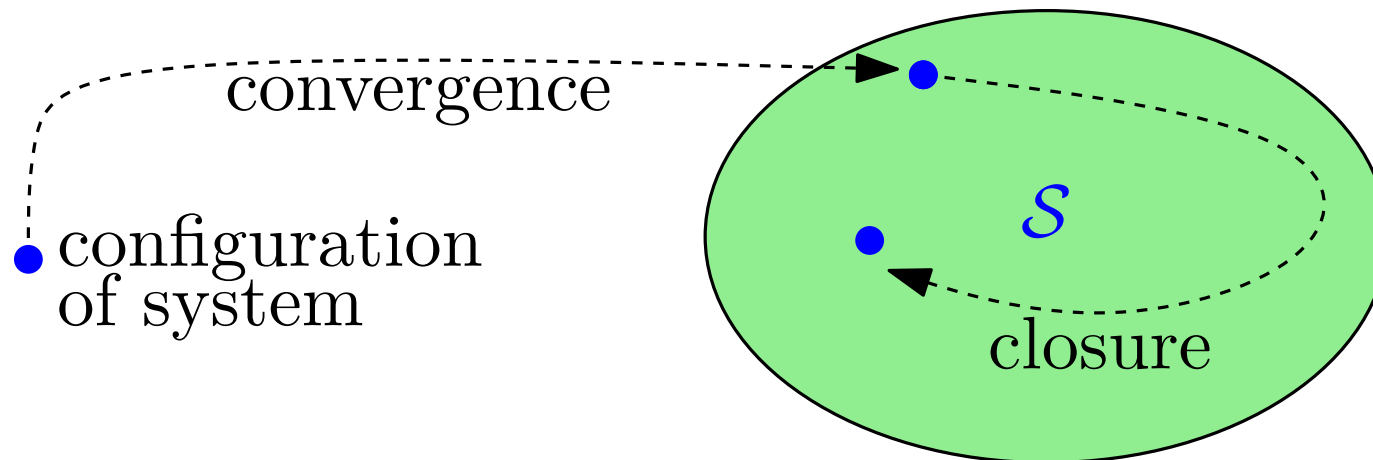
(Probabilistic) self-stabilization:

$\mathcal{S} := \{\text{“correct configurations of the system”}\}$
(= consensus on source's bit)

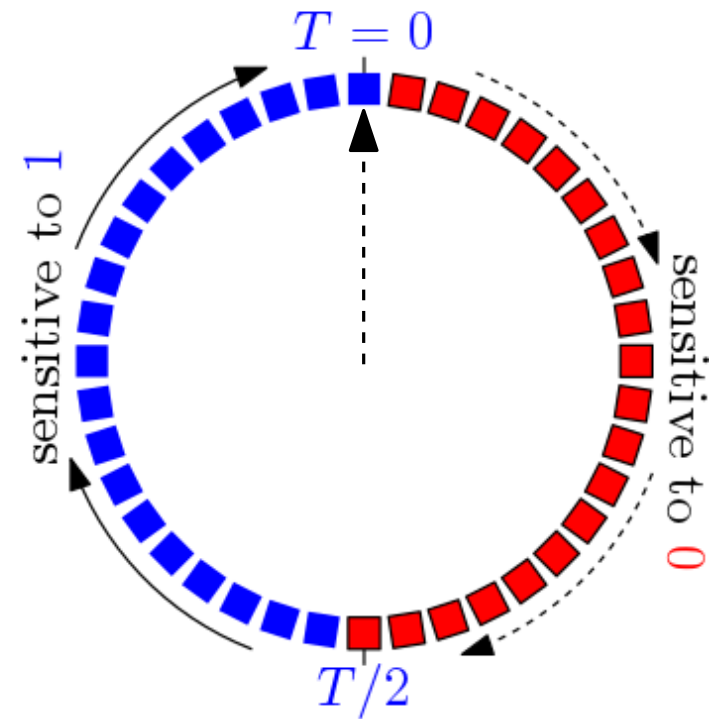
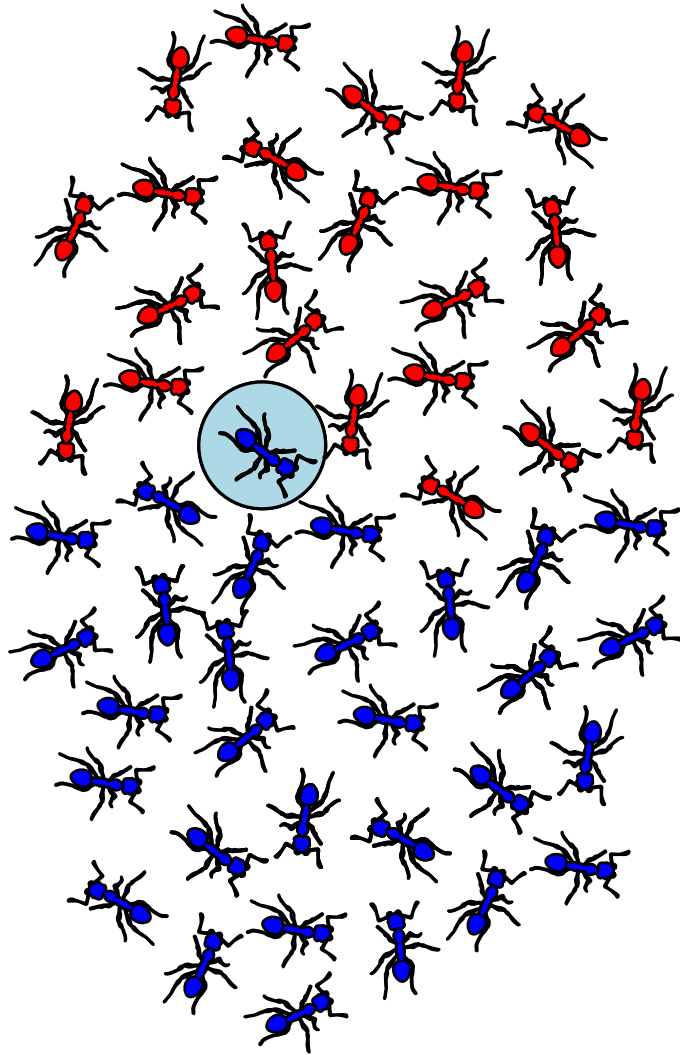
- **Convergence.** From *any* initial configuration, the system reaches \mathcal{S} (w.h.p.)
- **Closure.** If in \mathcal{S} , the system stays in \mathcal{S} (w.h.p.)

(Probabilistic) Self-stabilizing algorithm:

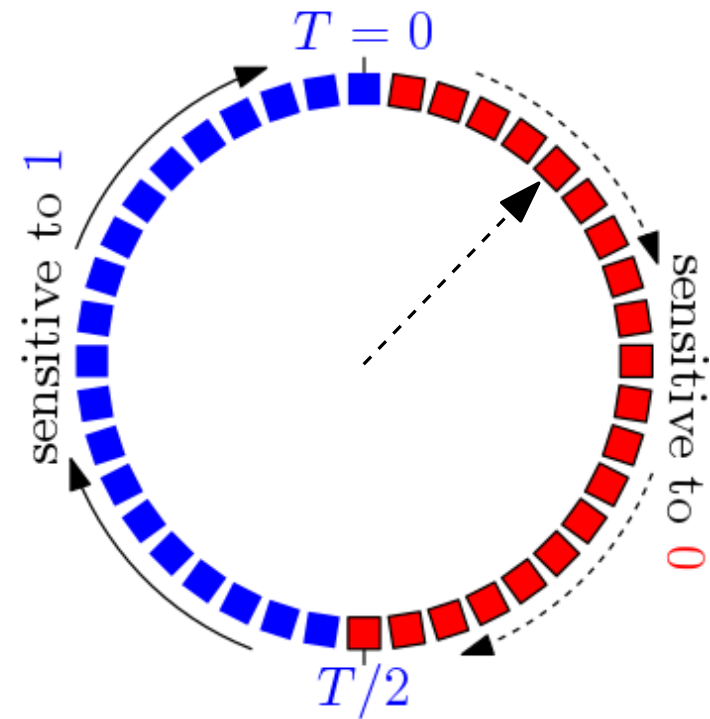
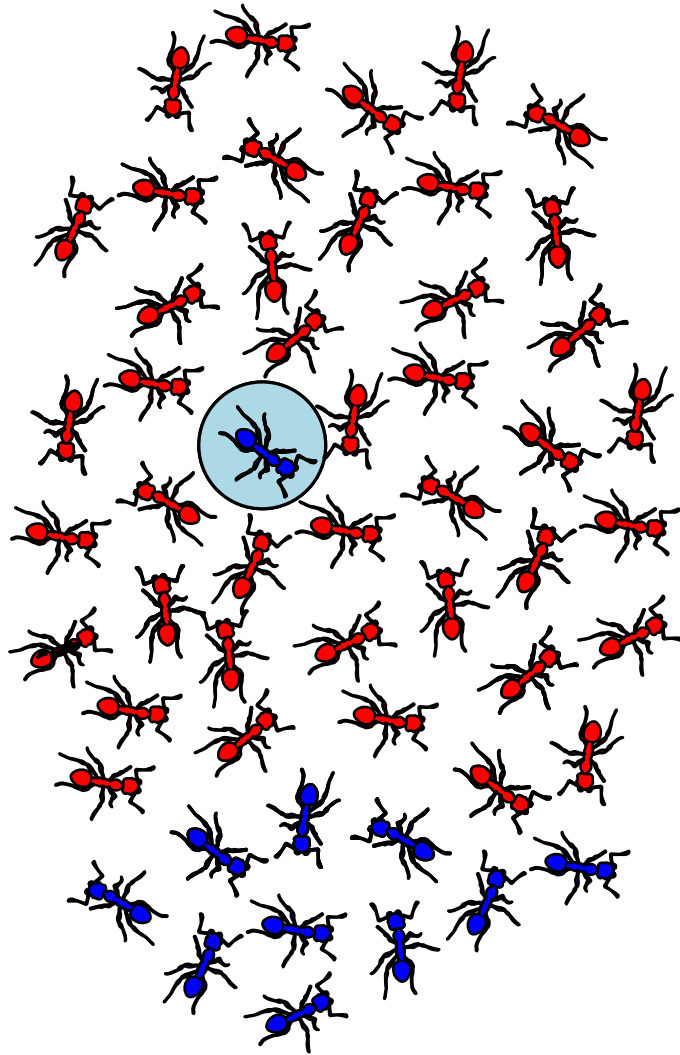
guarantees *convergence* and *closure* w.r.t. \mathcal{S} (*w.h.p.*)



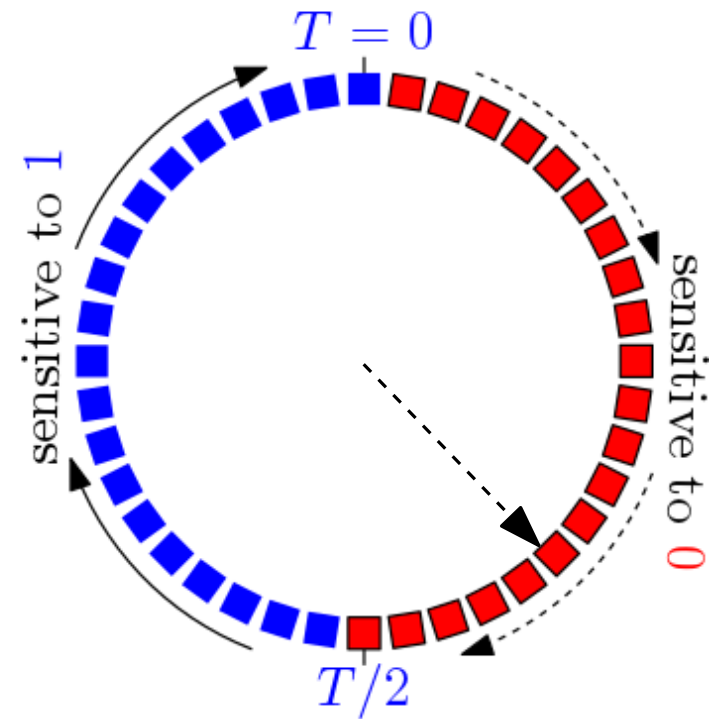
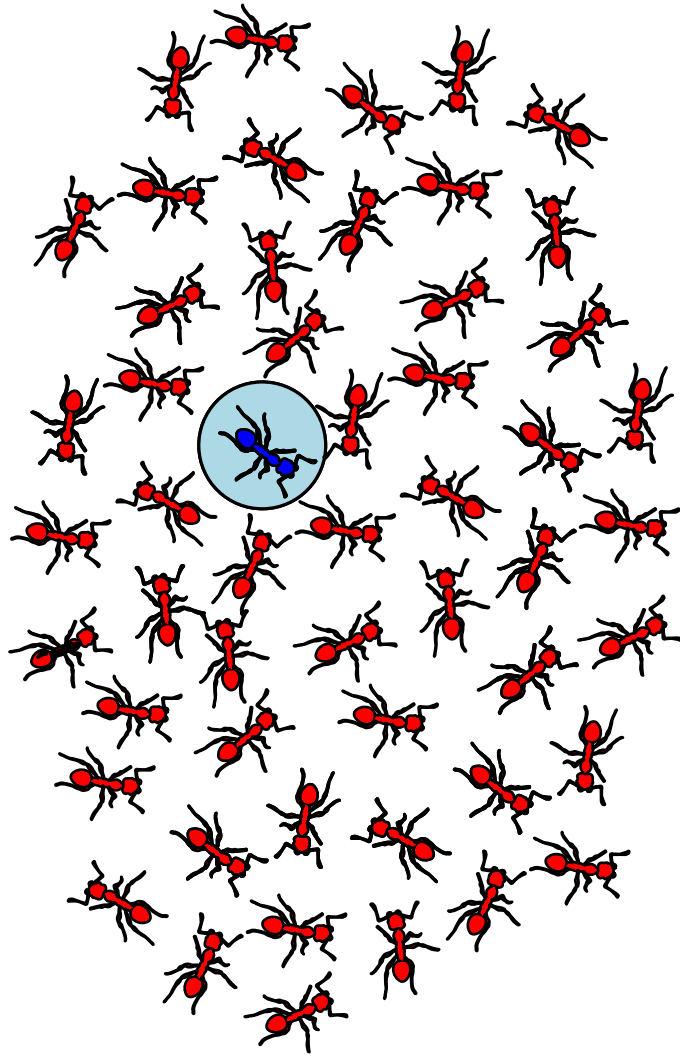
(Self-Stab.) Bit Dissemination vs Synchronization



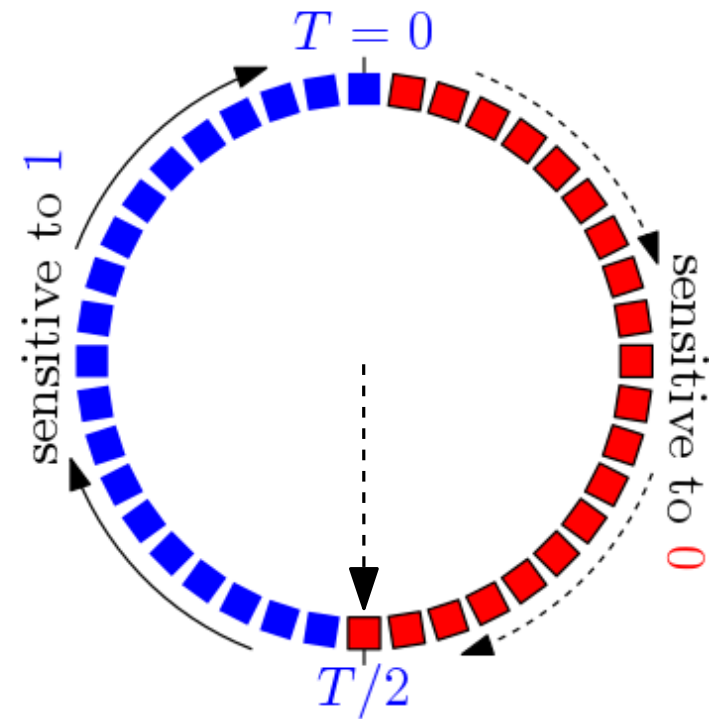
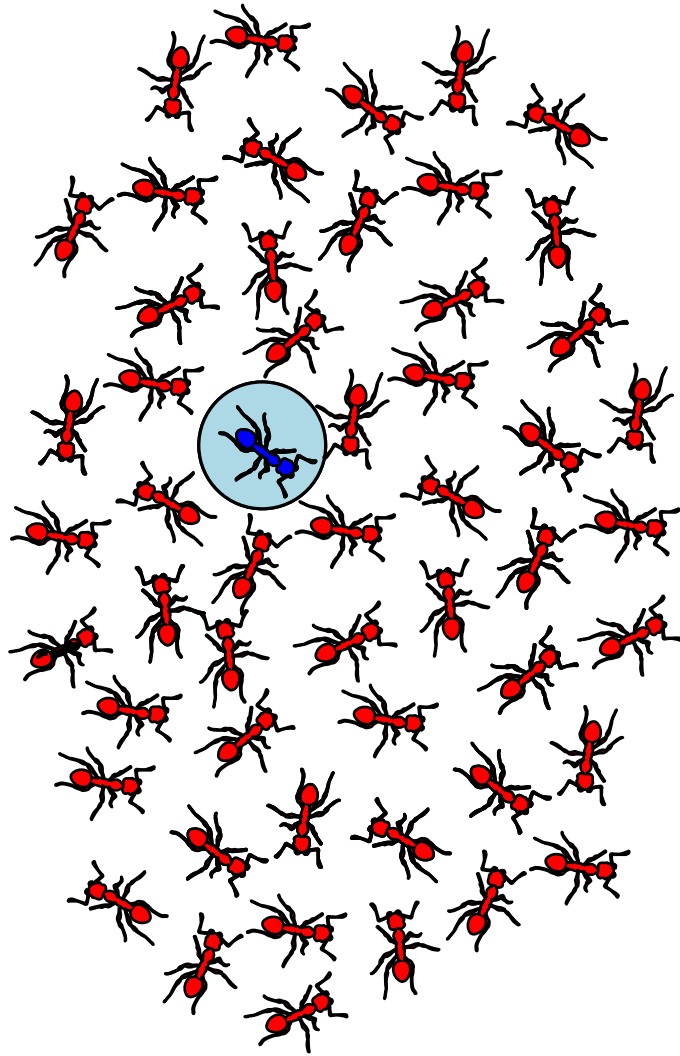
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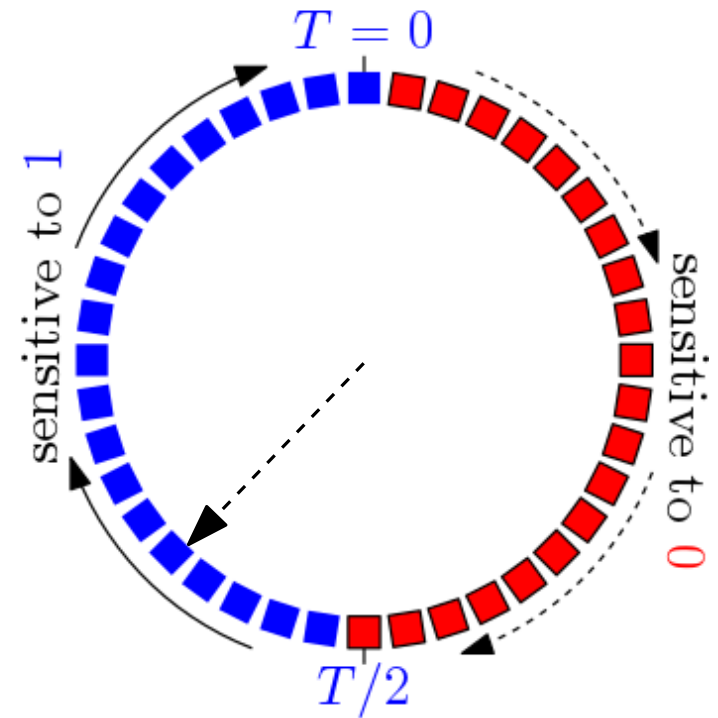
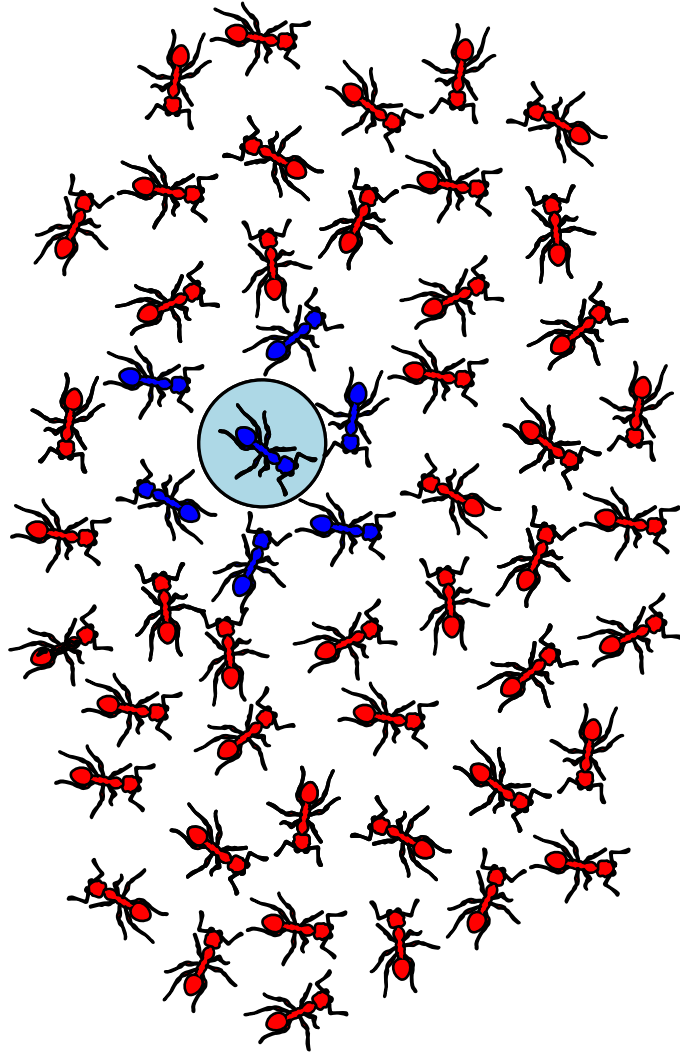
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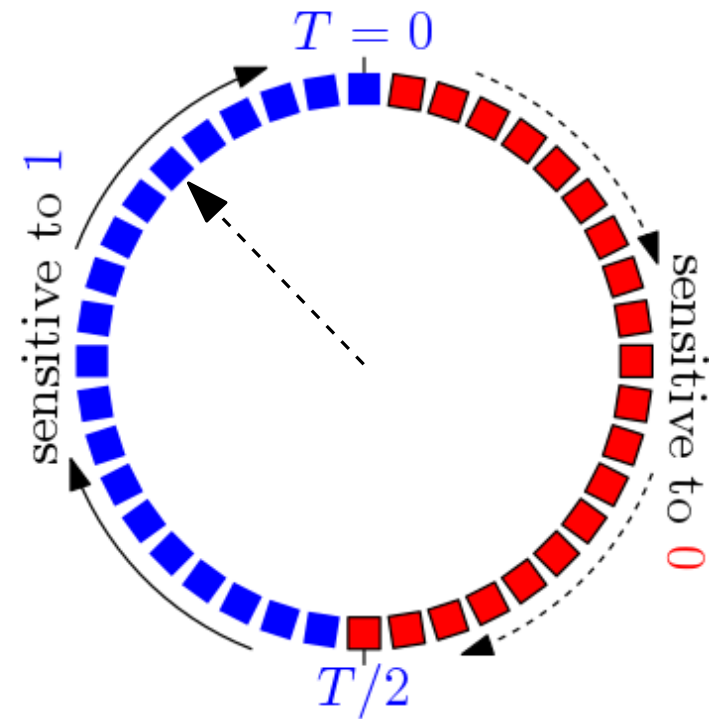
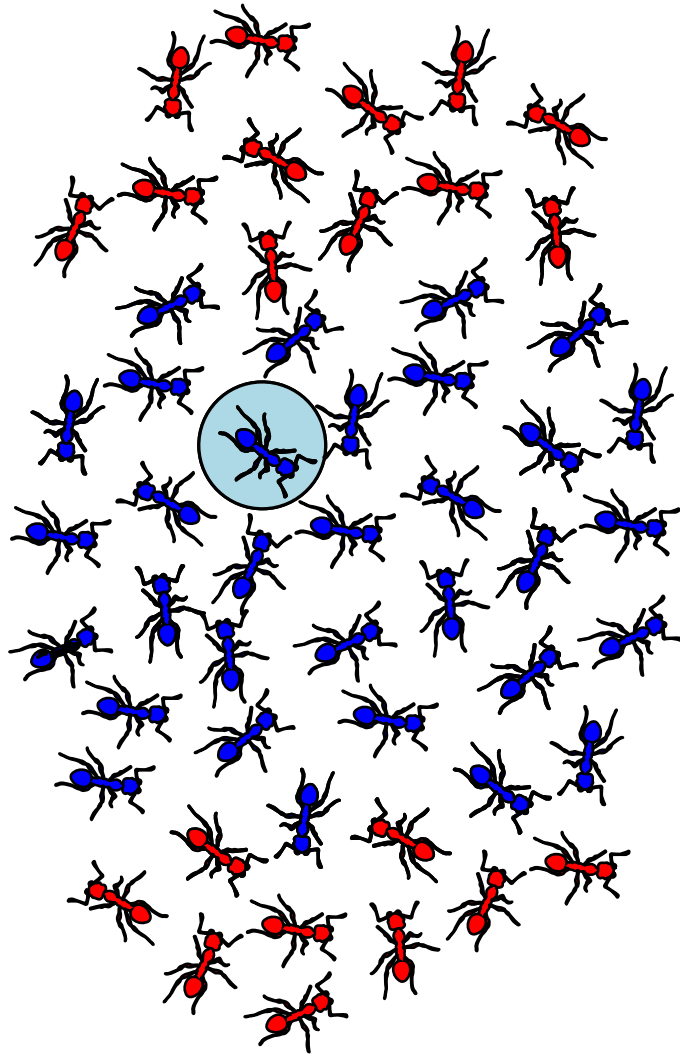
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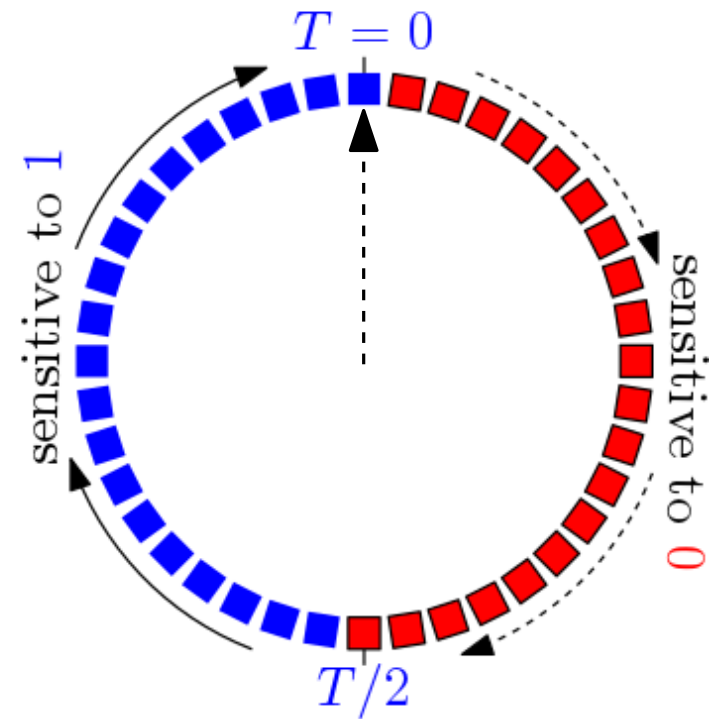
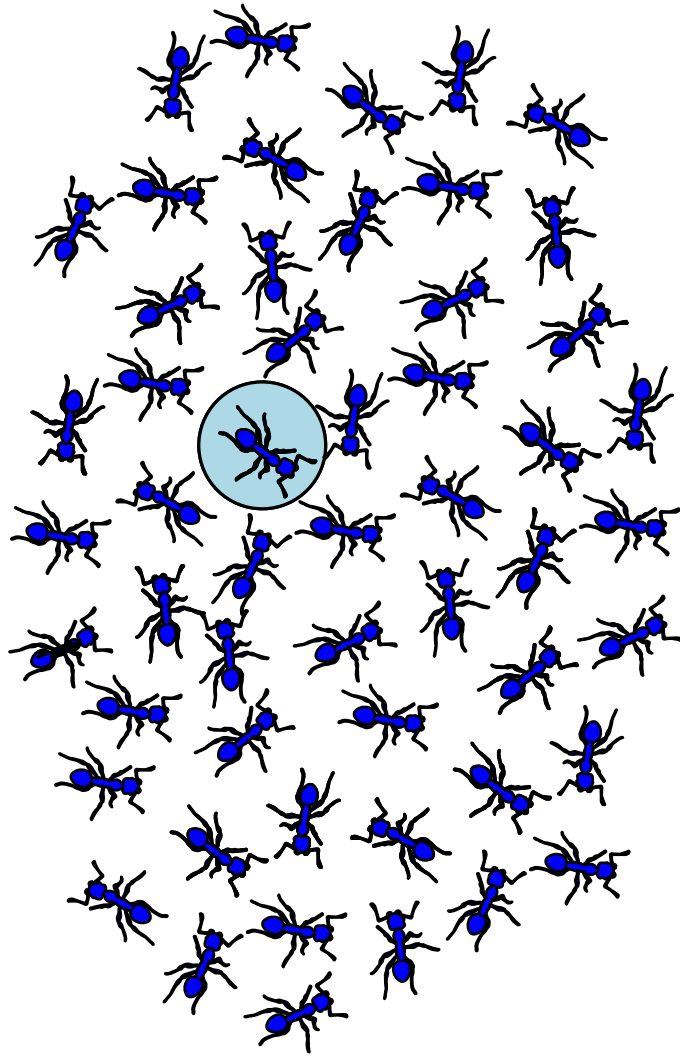
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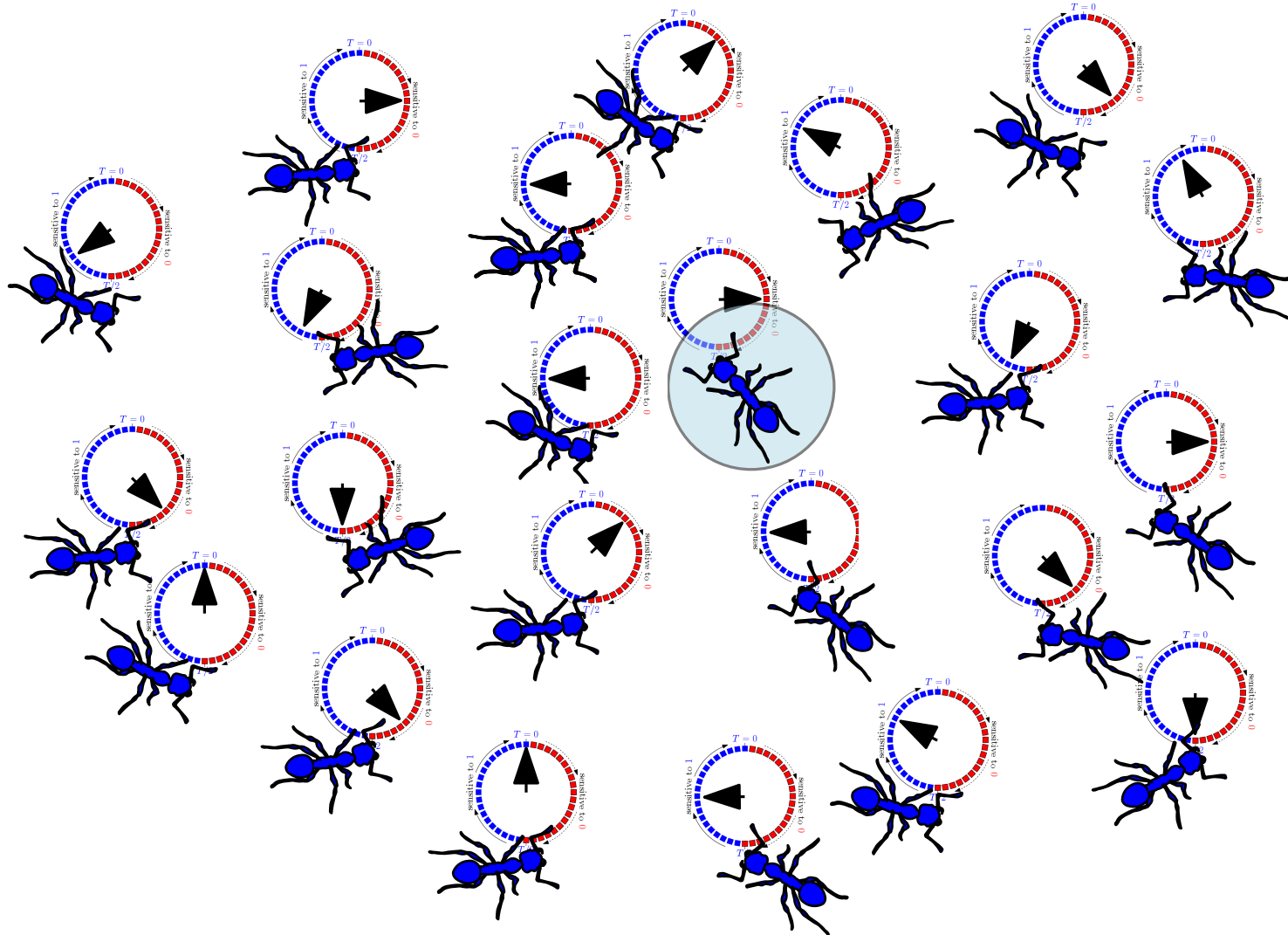


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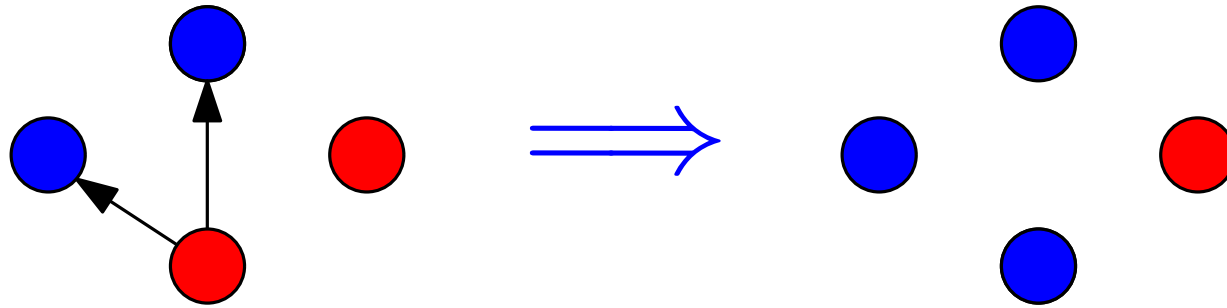


(Self-Stab.) Bit Dissemination vs Synchronization

Self-stabilizing algorithms converge from
any initial configuration

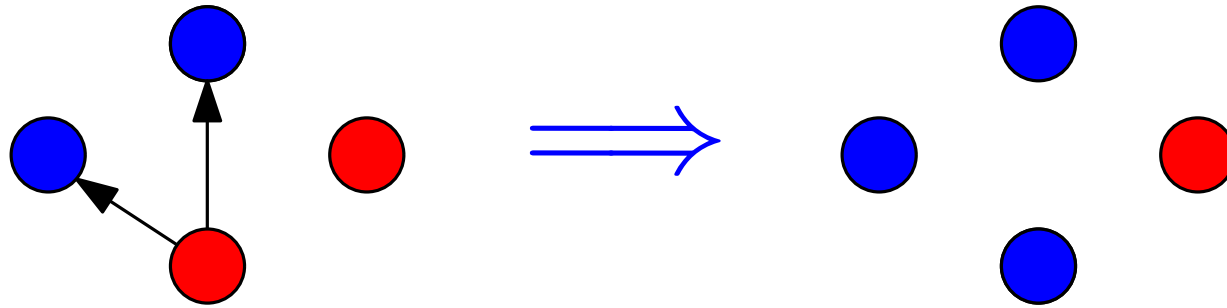


Self-Stabilizing Clock Sync. in the *PULL* Model

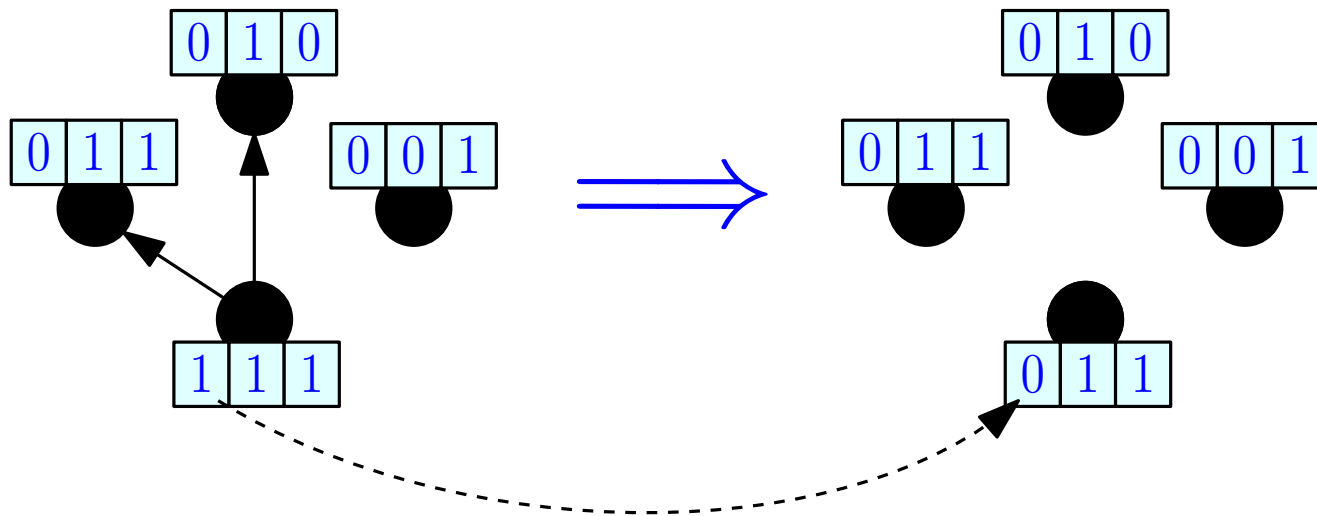


2-Majority dynamics [Doerr et al. '11]. Converge to consensus in $\mathcal{O}(\log n)$ rounds with high probability.

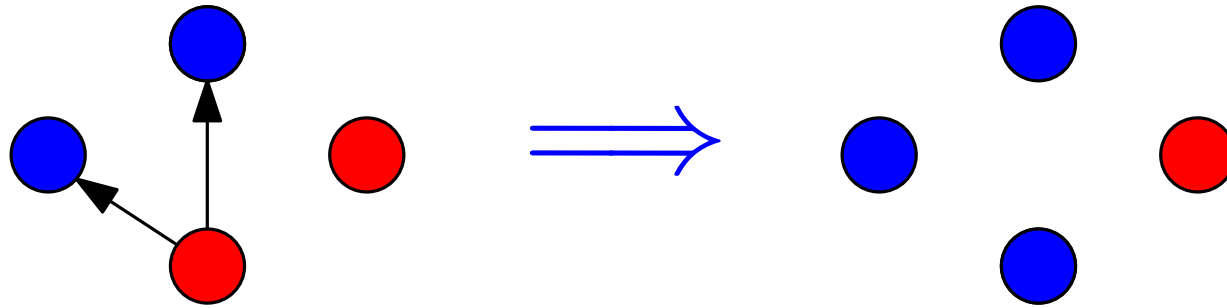
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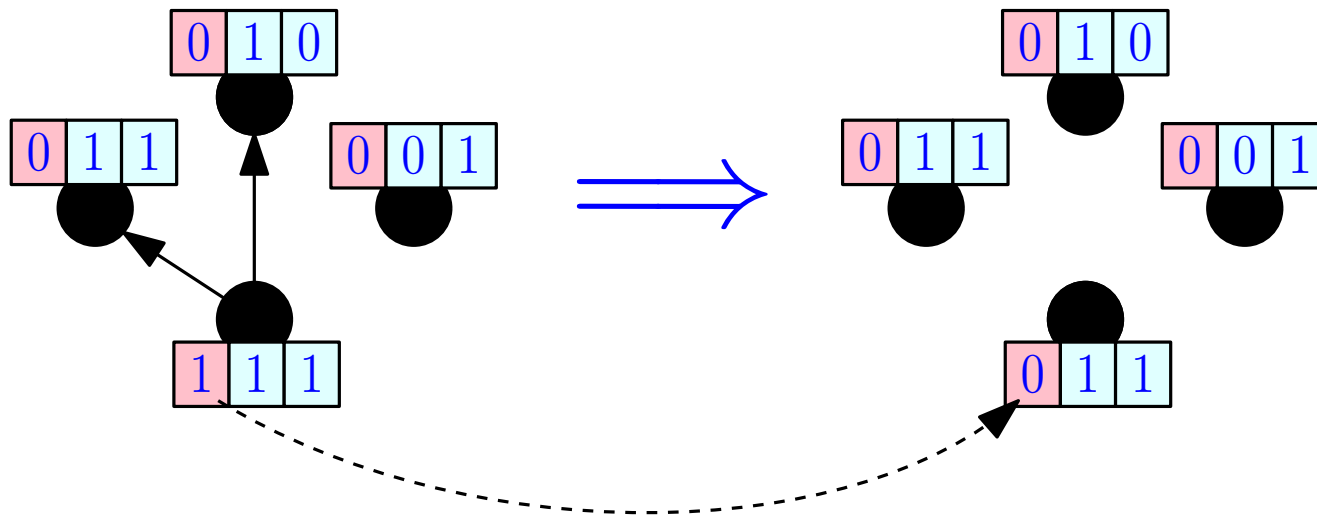
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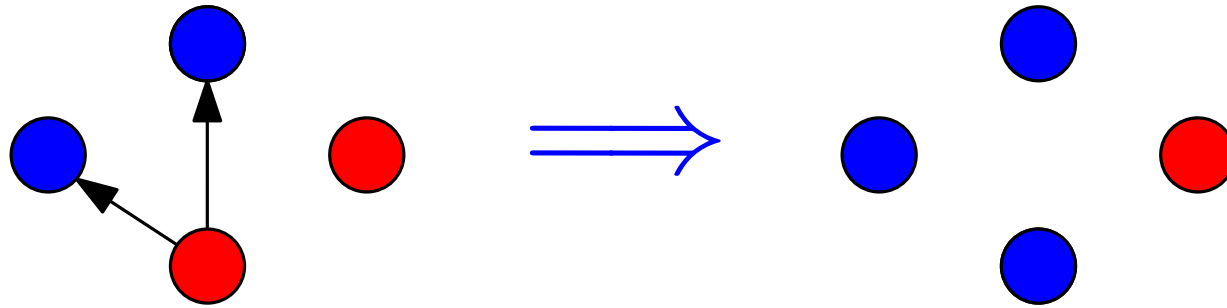
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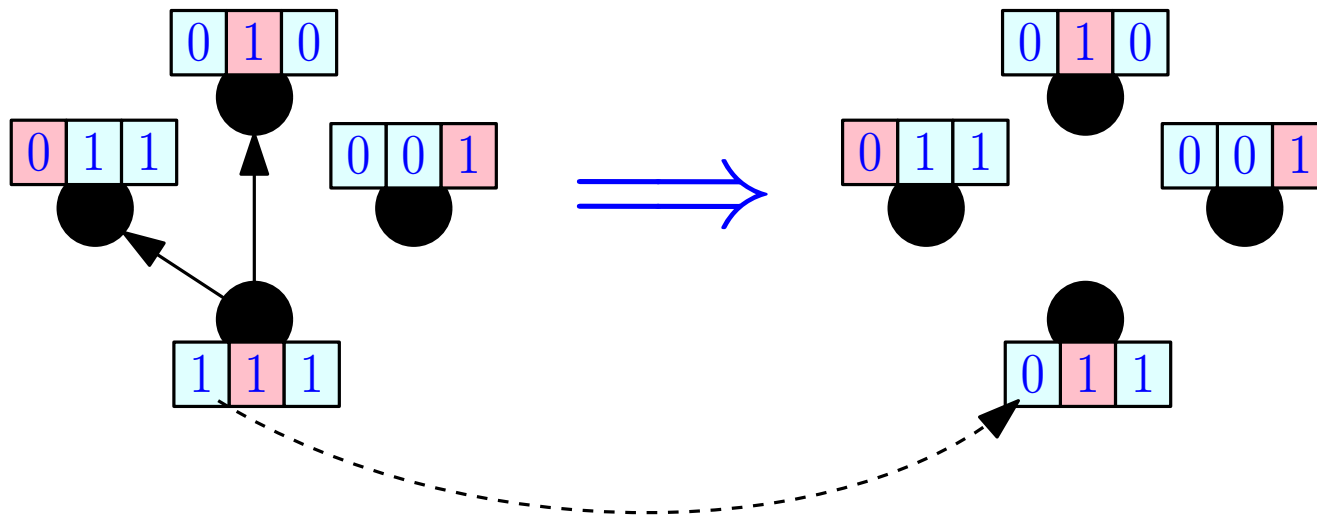
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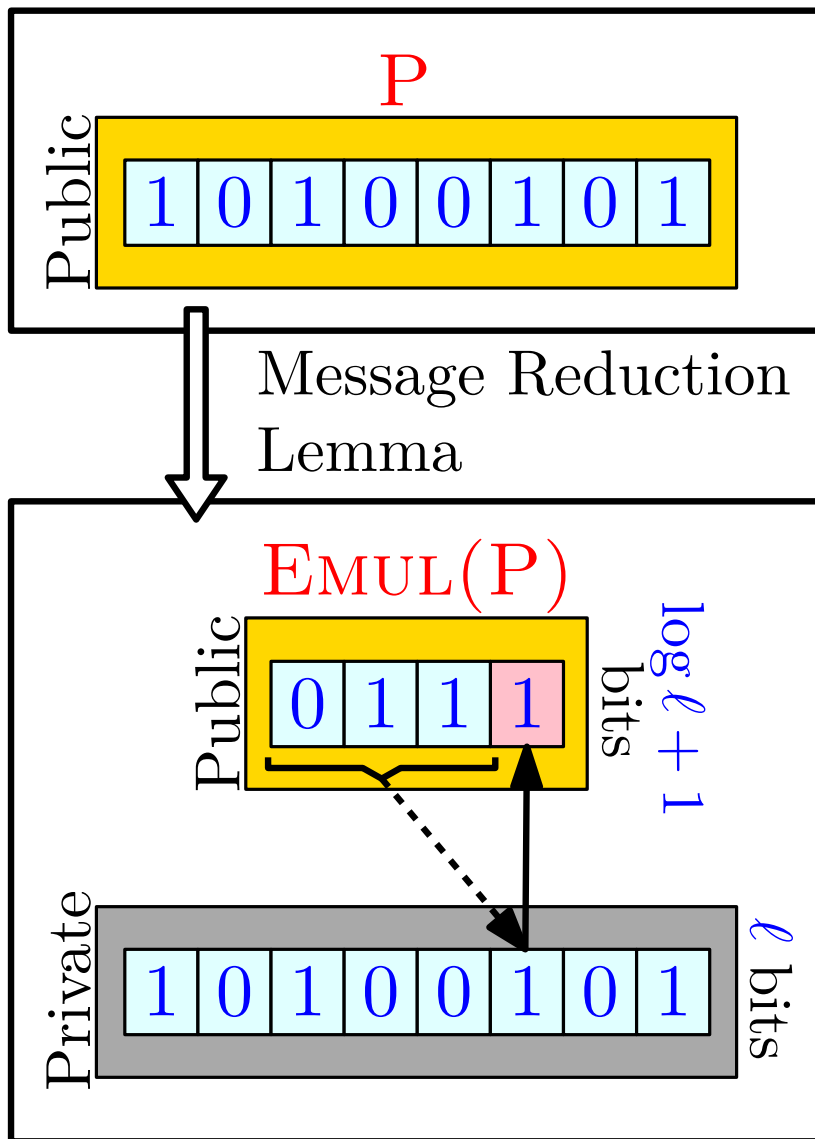
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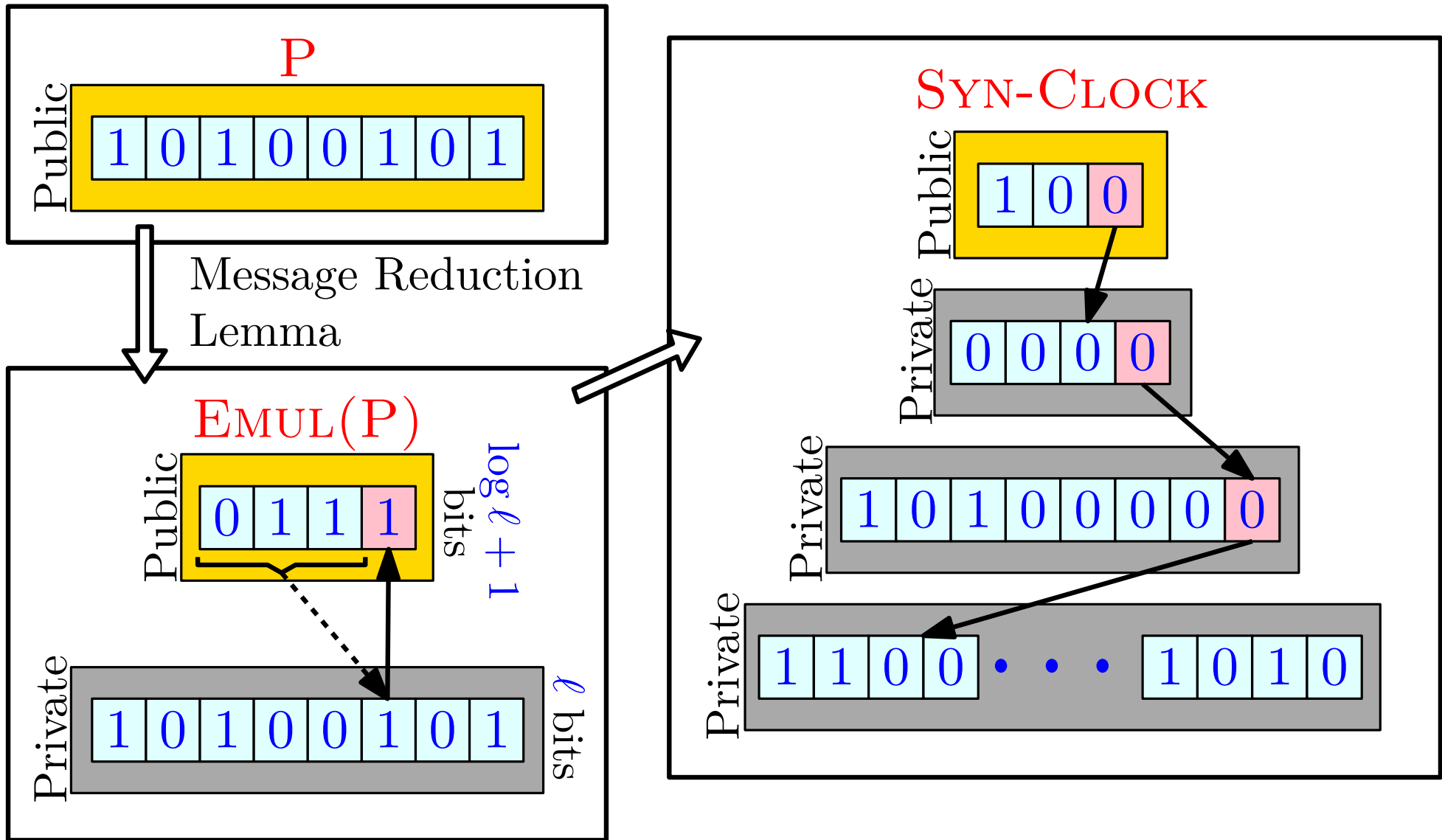
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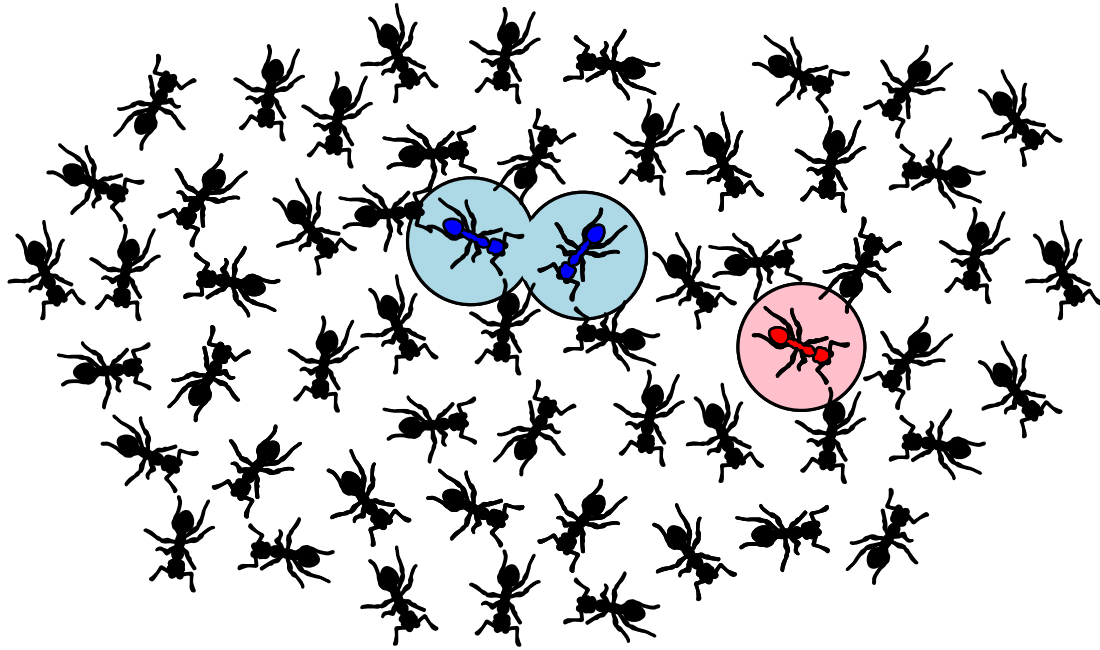
The Message Reduction Lemma



The Message Reduction Lemma



SYN-PHASE-SPREAD

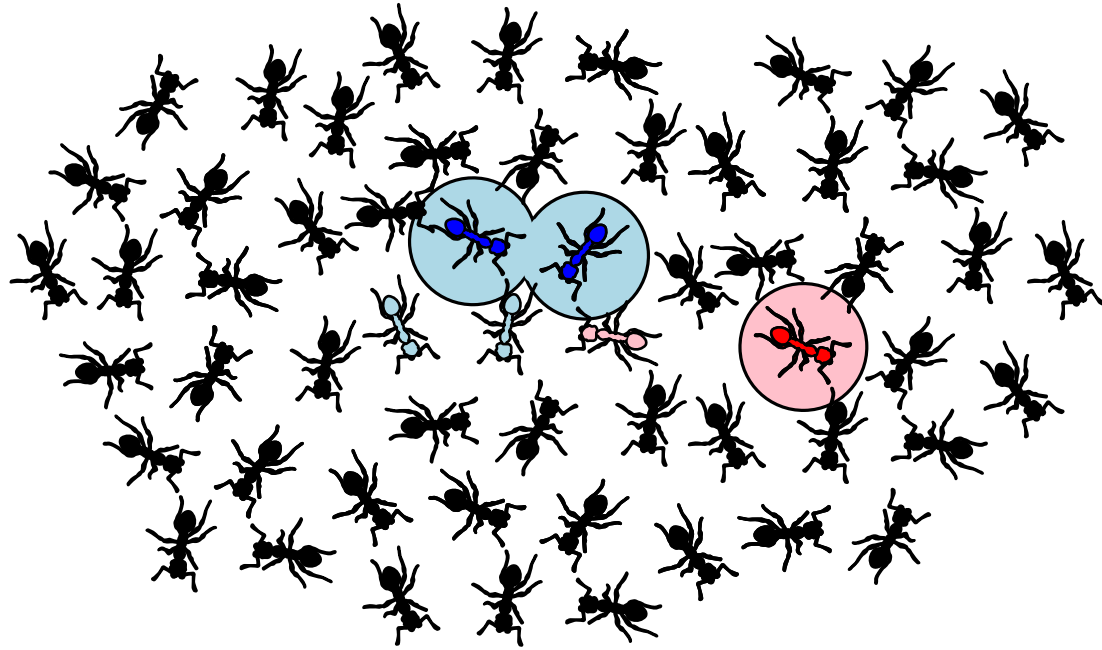


Spreading Phases
(Core idea: FHK'14)

blue vs red:

2/1

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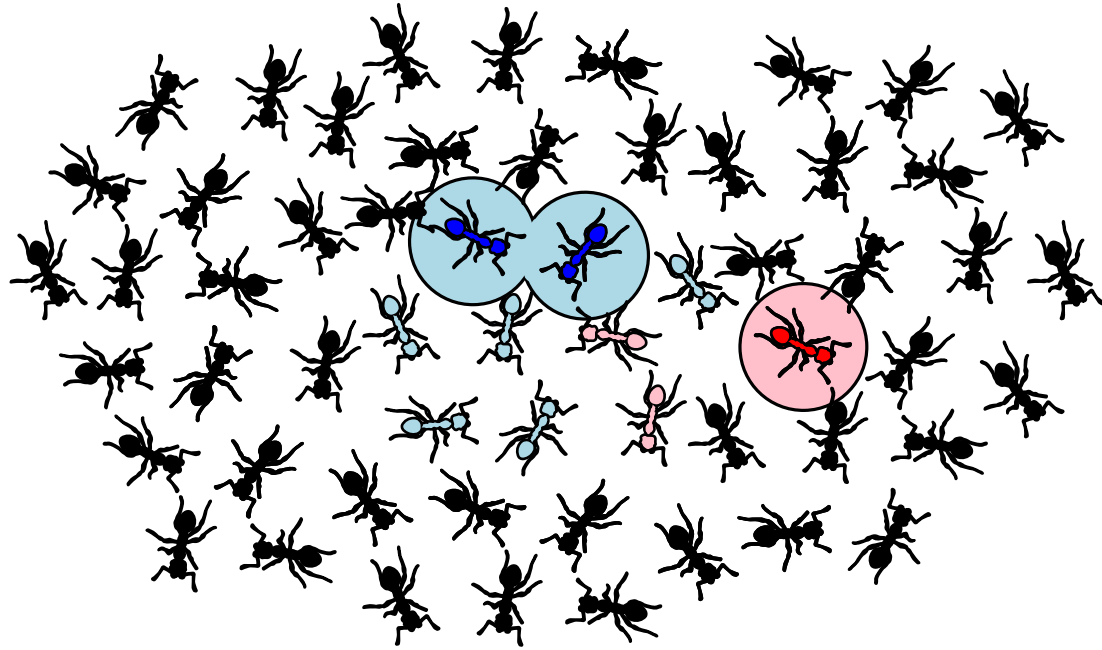


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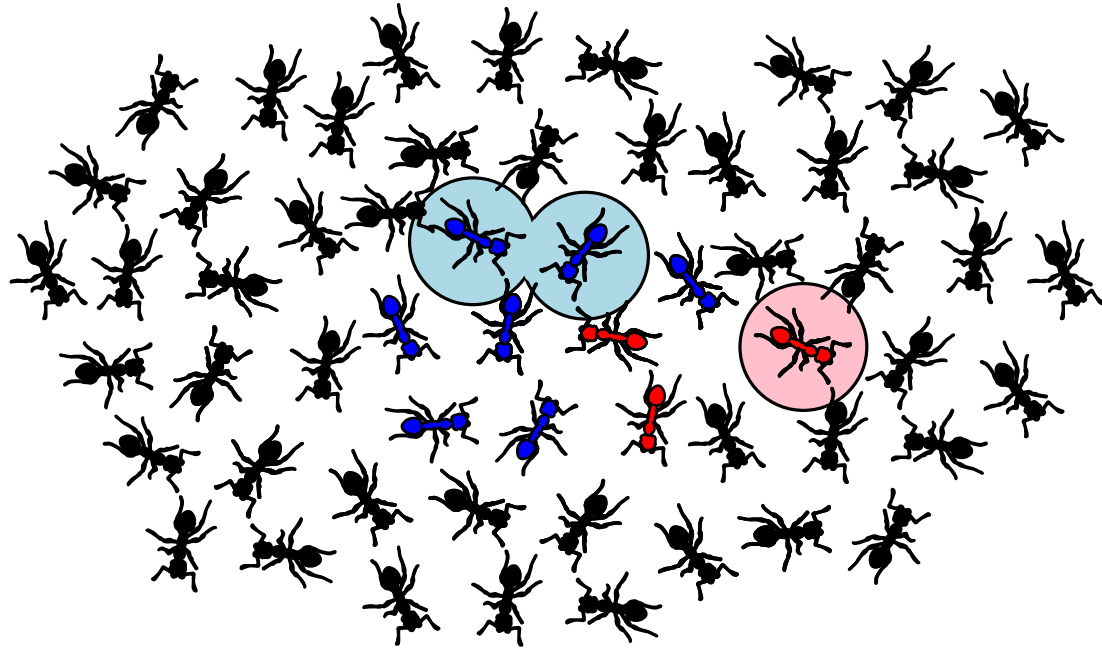


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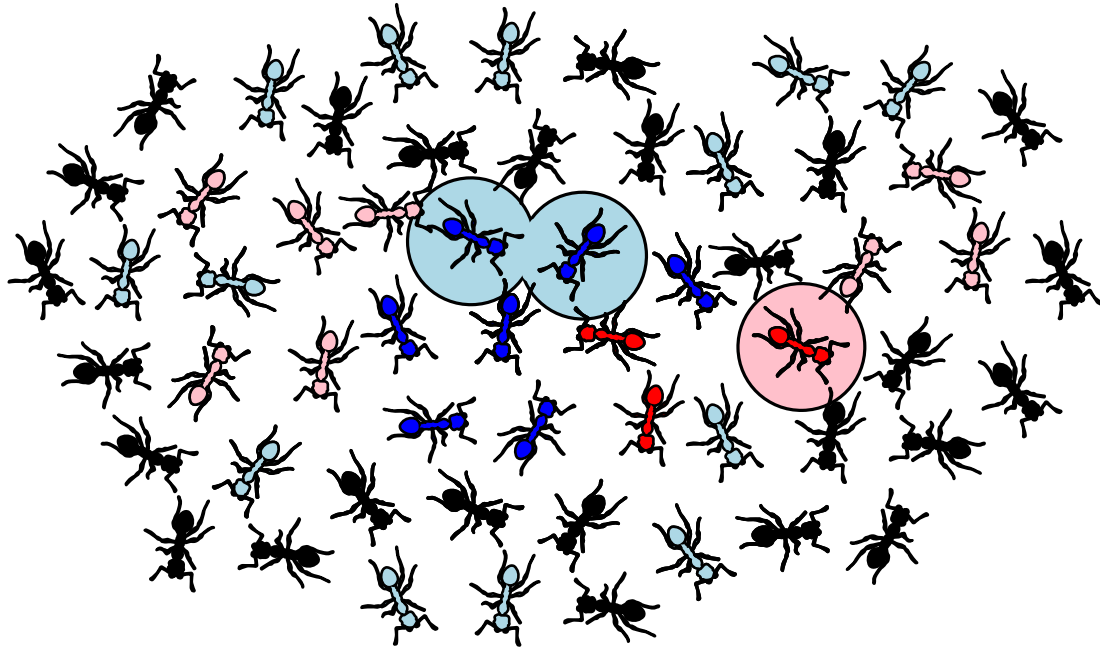


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8/4

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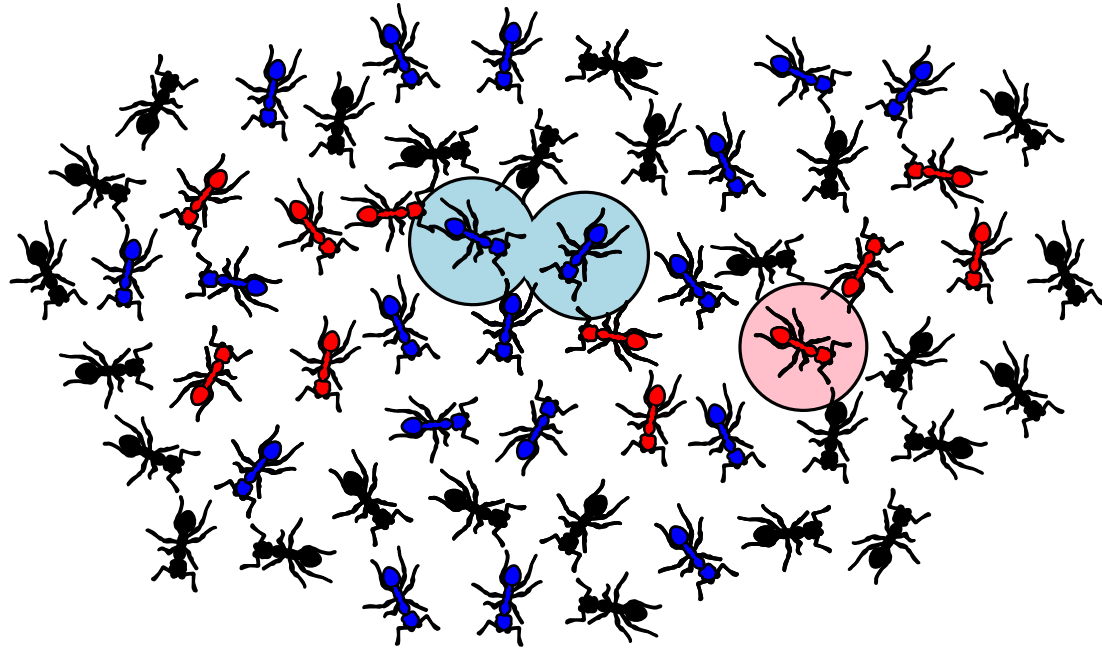


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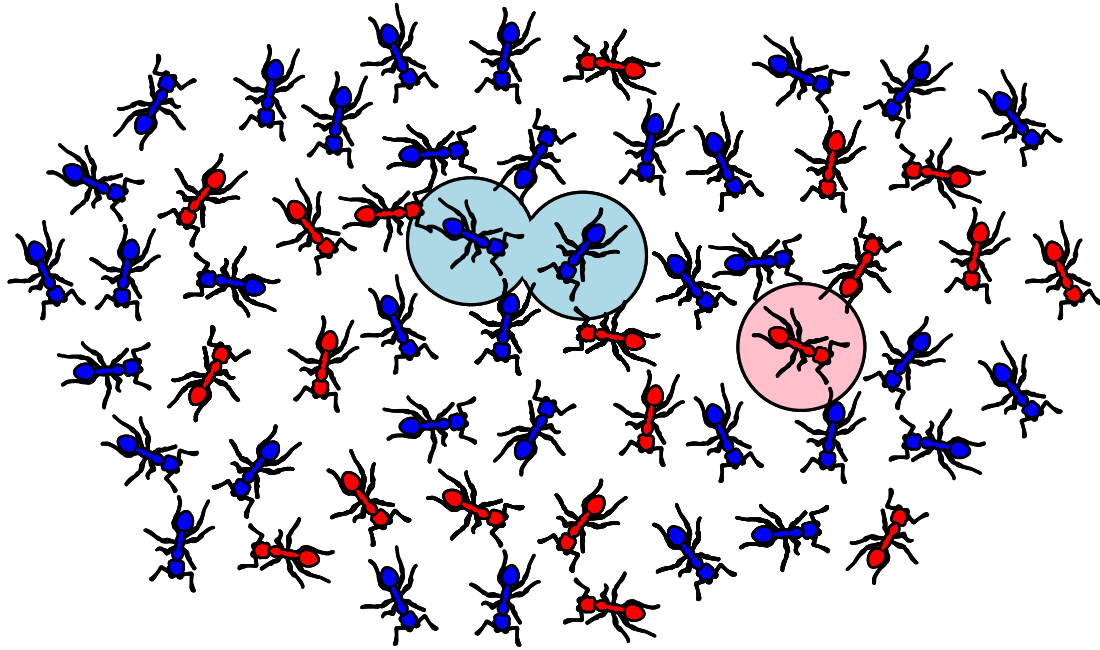


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$$20/11 \approx 1.85$$

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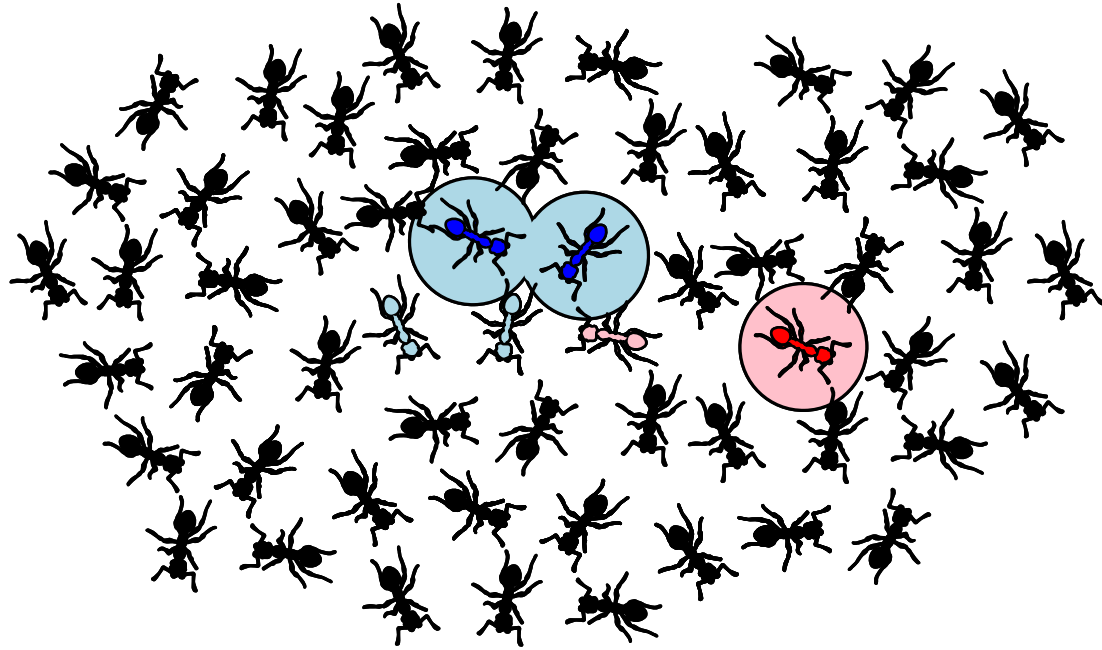


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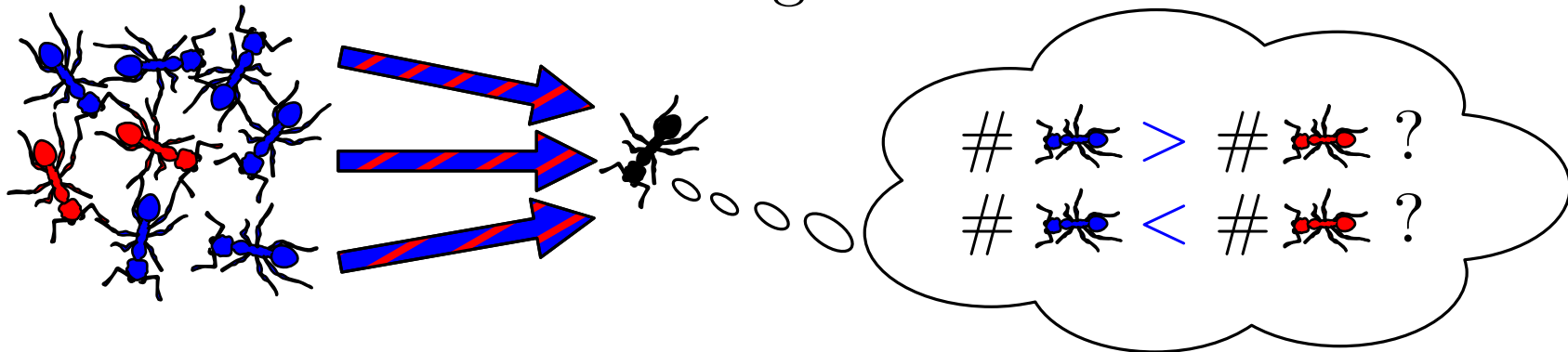
SYN-PHASE-SPREAD



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Polling Phase



Results

Theorem (Clock Synchronization). SYN-CLOCK is a *self-stabilizing* clock synchronization protocol which synchronizes a clock modulo T in $\tilde{O}(\log n \log T)$ rounds w.h.p. using 3-bit messages.

Theorem (Majority Bit Dissemination). SYN-PHASE-SPREAD is a *self-stabilizing* Majority Bit Dissemination protocol which converges in $\tilde{O}(\log n)$ rounds w.h.p using 3-bit messages, provided majority is supported by $(\frac{1}{2} + \epsilon)$ -fraction of source agents.

Self-Stab. Bit Diss. with 1 bit: a Candidate

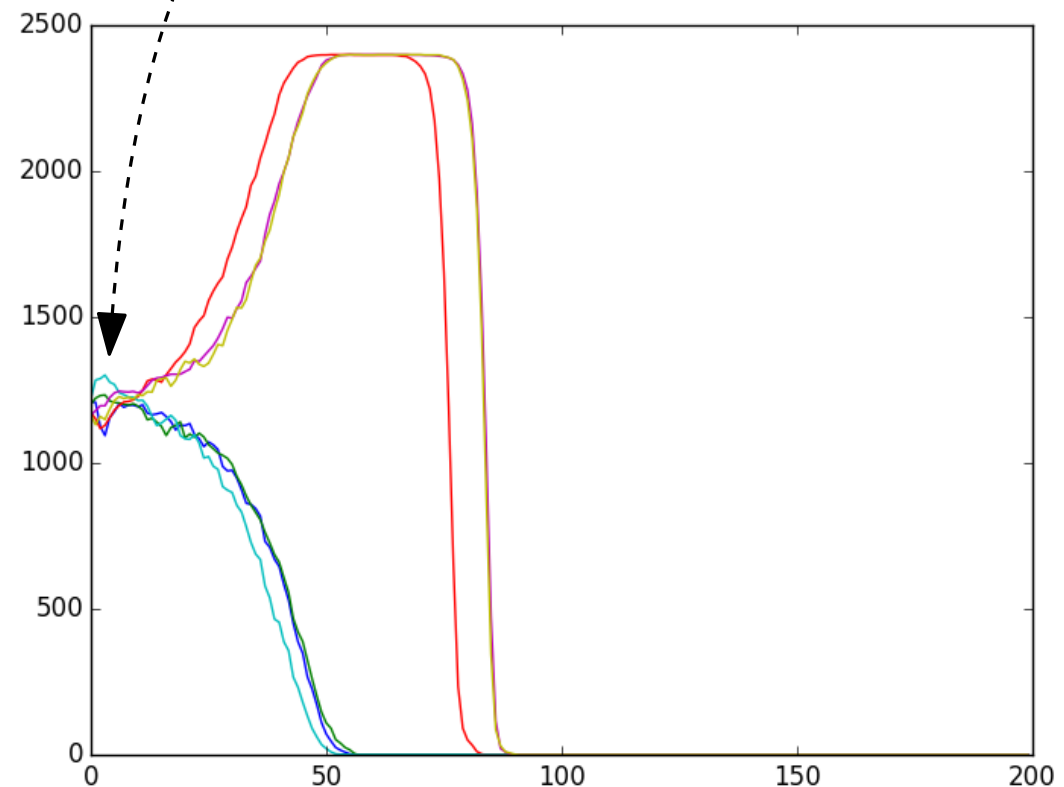
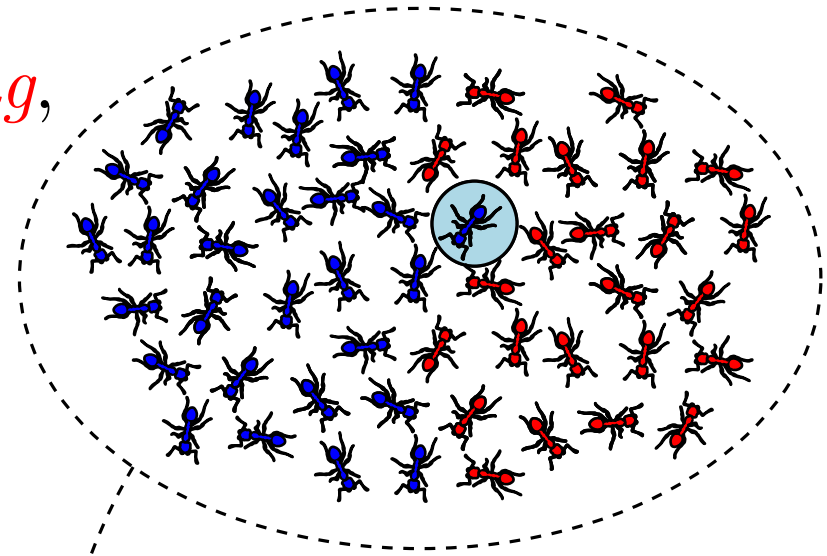
$\text{BFS}(f, s)$. Agents can *boosting*,
 $1/0$ -*frozen* or $1/0$ -*sensitive*.

- *Boosting*: Update their opinion with majority of their bit and the 2 bits they pull. If they see only agents of color c for s rounds, they become c -*sensitive*.
- c -*sensitive*: Turn into c -*frozen* if see value c .
- c -*frozen* keep value c for f rounds before becoming *boosting*.

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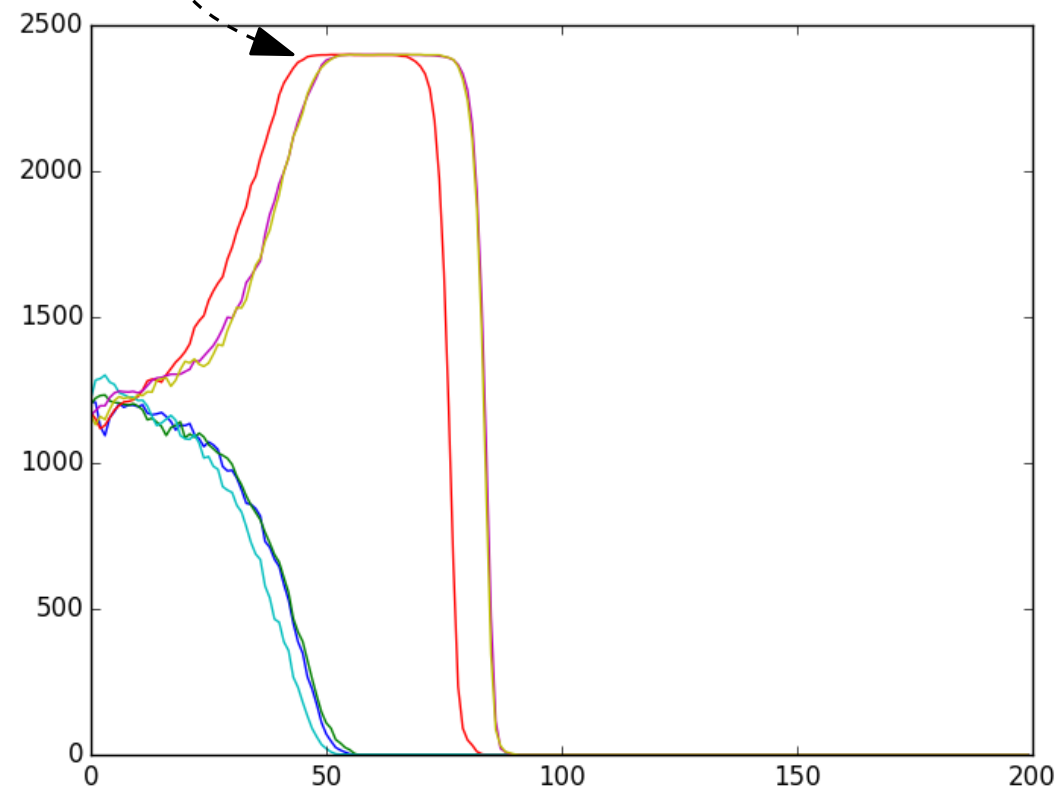
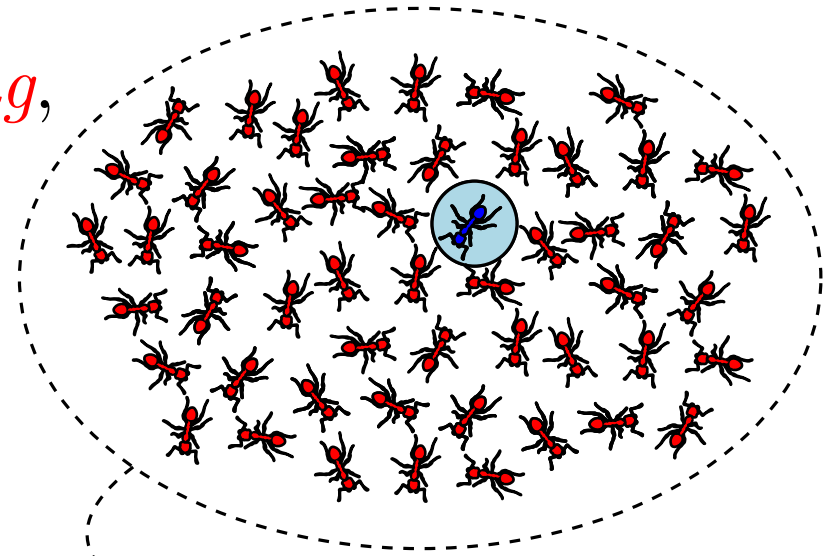
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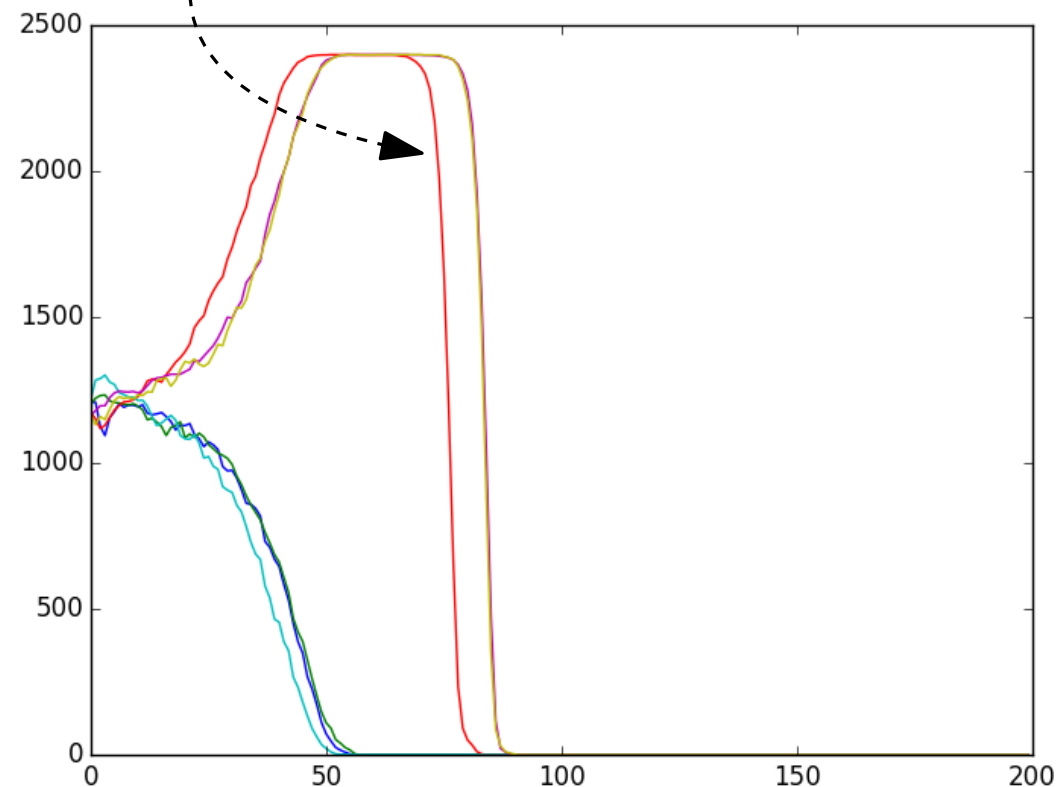
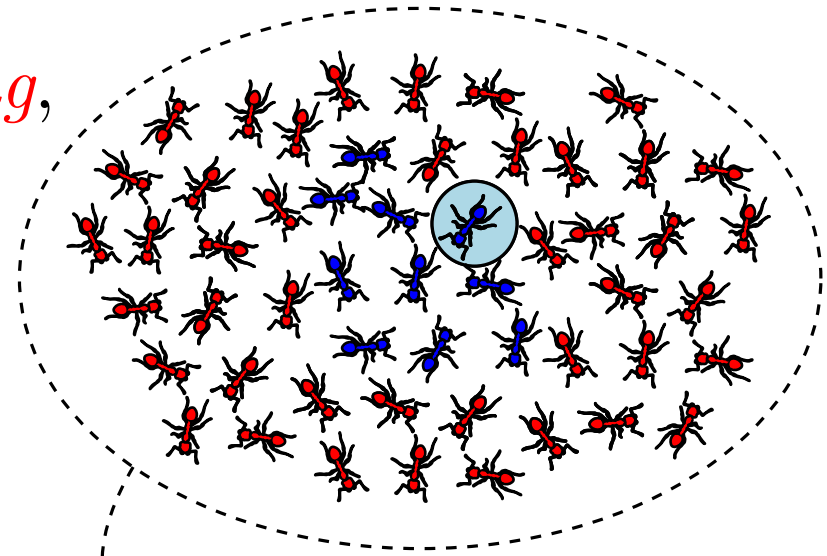
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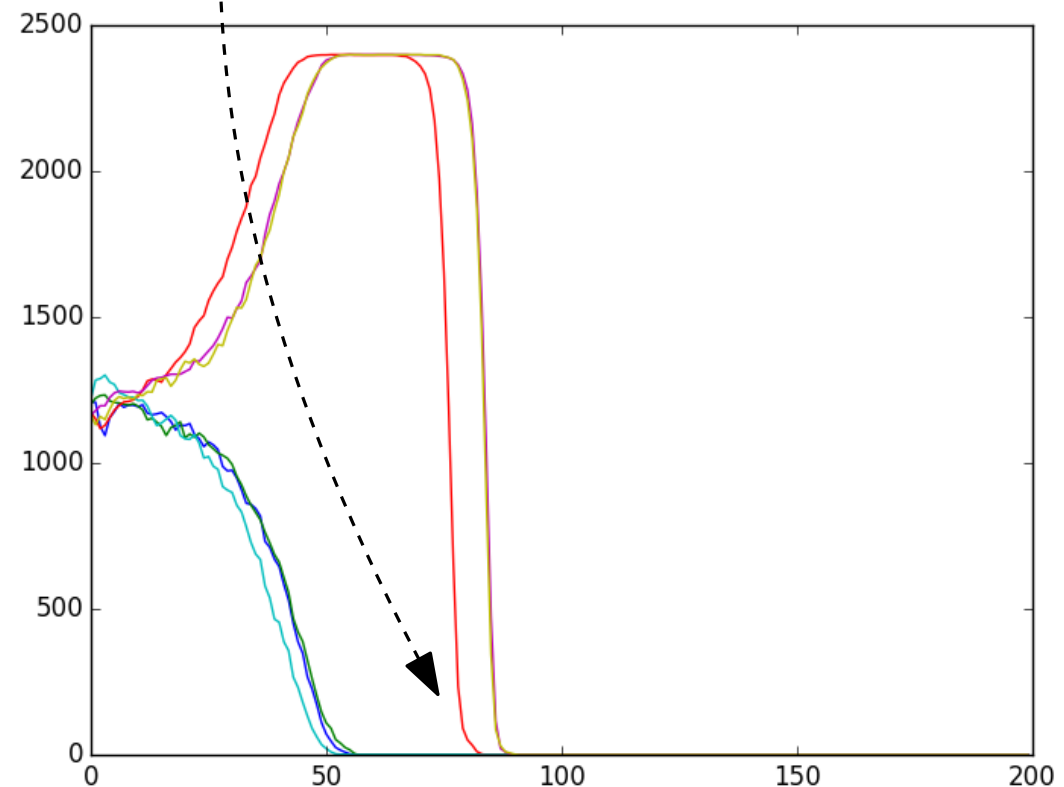
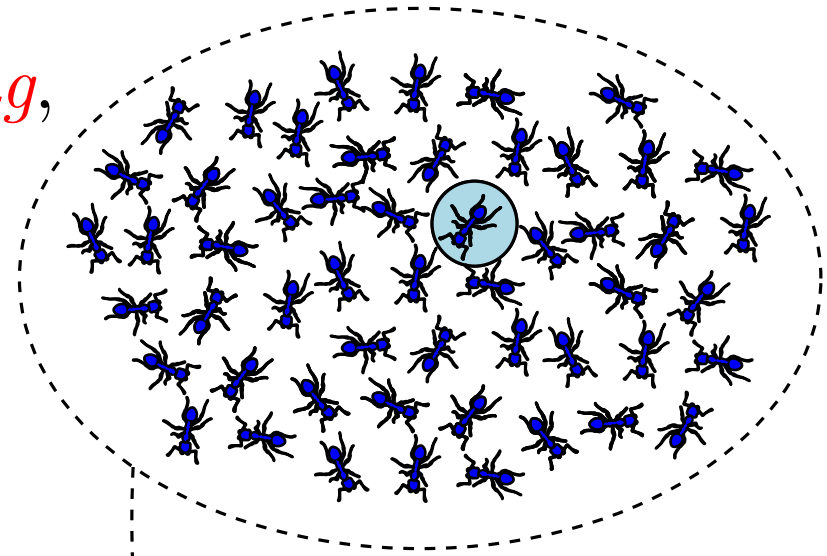
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- *Boosting*: Update their opinion with majority of their bit and the 2 bits they pull. If they see only agents of color c for s rounds, they become *c-sensitive*.
- *c-sensitive*: Turn into *c-frozen* if see value c .
- *c-frozen* keep value c for f rounds before becoming *boosting*.



Thank

You